

## When a partial Borel order is linearizable

by

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**Abstract.** We prove a classification theorem of the “Glimm–Effros” type for Borel order relations: a Borel partial order on the reals either is Borel linearizable or includes a copy of a certain Borel partial order  $\leq_0$  which is not Borel linearizable.

NOTATION. A binary relation  $\preceq$  on a set  $X$  is a *partial quasi-order*, or *p.q.-o.* in brief, on  $X$ , iff  $x \preceq y \wedge y \preceq z \Rightarrow x \preceq z$ , and  $x \preceq x$  for any  $x \in X$ . In this case,  $\approx$  is the associated equivalence relation, i.e.  $x \approx y$  iff  $x \preceq y \wedge y \preceq x$ .

If in addition  $x \approx x \Rightarrow x = x$  for any  $x$  then  $\preceq$  is a *partial order*, or *p.o.*, so that, say, forcing relations are p.q.-o.’s, but, generally speaking, not p.o.’s in this terminology.

A p.o. is *linear* (l.o.) iff we have  $x \preceq y \vee y \preceq x$  for all  $x, y \in X$ .

Let  $\preceq$  and  $\preceq'$  be p.q.-o.’s on resp.  $X$  and  $X'$ . A map  $h : X \rightarrow X'$  will be called *half order preserving*, or *h.o.p.*, iff  $x \preceq y \Rightarrow h(x) \preceq' h(y)$ .

DEFINITION 1. A Borel p.q.-o.  $\langle X; \preceq \rangle$  is *Borel linearizable* iff there is a Borel l.o.  $\langle X'; \preceq' \rangle$  and a Borel h.o.p. map  $h : X \rightarrow X'$  (called a *linearization map*) satisfying  $x \approx y \Leftrightarrow h(x) = h(y)$  <sup>(1)</sup>.

**Introduction.** Harrington, Marker, and Shelah [2] proved several theorems on Borel partial order relations, mainly concerning *thin* p.q.-o.’s, i.e. those which do not admit uncountable pairwise incomparable subsets. In particular, they demonstrated that any such Borel p.q.-o. is Borel linearizable, and moreover the corresponding l.o.  $\langle X'; \preceq' \rangle$  can be chosen as a suborder of  $\langle 2^\alpha; \leq_{\text{lex}} \rangle$  for some  $\alpha < \omega_1$ , where  $\leq_{\text{lex}}$  is the lexicographical order.

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<sup>(1)</sup> The equivalence cannot be dropped in this definition as otherwise a one-element set  $X'$  works in any case.

As elementary examples show that thinness is not a necessary condition for Borel linearizability, this result leaves open the problem of linearization of non-thin Borel p.q.-o.'s. Harrington *et al.* wrote in [2] that “there is little to say about nonthin orderings”, although there are many interesting among them like the *dominance* order on  $\omega^\omega$ .

Our main result will say that not all Borel p.q.-o.'s are Borel linearizable, and there exists a *minimal* one, in a certain sense, among them.

DEFINITION 2. Let  $a, b \in 2^\omega$ . We define  $a \leq_0 b$  iff either  $a = b$  or  $a E_0 b$  <sup>(2)</sup> and  $a(k_0) < b(k_0)$  where  $k_0$  is the largest  $k$  such that  $a(k) \neq b(k)$  <sup>(3)</sup>.

The relation  $\leq_0$  is a Borel p.q.-o. on  $2^\omega$  which orders every  $E_0$ -class similarly to the integers  $\mathbb{Z}$  (except for the class  $[\omega \times \{0\}]_{E_0}$  ordered as  $\omega$  and the class  $[\omega \times \{1\}]_{E_0}$  ordered as  $\omega^*$ , the inverted order) but leaves any two  $E_0$ -inequivalent reals incomparable.

The following is the main result of the paper.

THEOREM 3. Suppose that  $\preceq$  is a Borel p.q.-o. on  $\mathcal{N} = \omega^\omega$ . Then exactly one of the following two conditions is satisfied:

(I)  $\preceq$  is Borel linearizable; moreover <sup>(4)</sup>, there exist an ordinal  $\alpha < \omega_1$  and a Borel linearization map  $h : \langle \mathcal{N}; \preceq \rangle \rightarrow \langle 2^\alpha; \leq_{\text{lex}} \rangle$ .

(II) there exists a continuous 1-1 map  $F : 2^\omega \rightarrow \mathcal{N}$  such that we have  $a \leq_0 b \Rightarrow F(a) \preceq F(b)$  while  $a E_0 b$  implies that  $F(a)$  and  $F(b)$  are  $\preceq$ -incomparable <sup>(5)</sup>.

The theorem resembles the case of Borel equivalence relations where a necessary and sufficient condition for a Borel equivalence relation  $E$  to be *smooth* is that  $E_0$  (which is not smooth) does not continuously embed in  $E$  (Harrington, Kechris and Louveau [1]). ( $\leq_0$  itself is *not* Borel linearizable.)

The proof is essentially a combination of ideas and techniques in [1, 2].

**1. Incompatibility.** Let us first prove that (I) and (II) are incompatible.

<sup>(2)</sup> That is,  $a(k) = b(k)$  for all but finite  $k$ , the *Vitali* equivalence relation on  $2^\omega$ .

<sup>(3)</sup> If one enlarges  $\leq_0$  so that, in addition,  $a <_0 b$  whenever  $a, b \in 2^\omega$  are such that  $a(k) = 1$  and  $b(k) = 0$  for all but finite  $k$  then the enlarged relation can be induced by a Borel action of  $\mathbb{Z}$  on  $2^\omega$ , such that  $a <_0 b$  iff  $a = zb$  for some  $z \in \mathbb{Z}, z > 0$ .

<sup>(4)</sup> The “moreover” assertion is an immediate corollary of the linearizability by the above-mentioned result of [2].

<sup>(5)</sup> Then  $F$  associates a chain  $\{F(b) : b E_0 a\}$  in  $\langle \mathcal{N}; \preceq \rangle$  to each  $E_0$ -class  $[a]_{E_0}$  so that any two different chains do not contain  $\preceq$  comparable elements: let us call them *fully incomparable* chains. Thus (II) essentially says that  $\preceq$  admits an effectively “large” Borel family of fully incomparable chains, which is therefore necessary and sufficient for  $\preceq$  to be *not* Borel linearizable.

Suppose otherwise. The superposition of the maps  $F$  and  $h$  is then a Borel h.o.p. map  $\phi : \langle 2^\omega; \leq_0 \rangle \rightarrow \langle 2^\alpha; \leq_{\text{lex}} \rangle$  satisfying the following:  $\phi(a) = \phi(b)$  implies that  $a E_0 b$ , i.e.  $a$  and  $b$  are  $\leq_0$  comparable.

Therefore, as any  $E_0$ -class is  $\leq_0$ -ordered similarly to  $\mathbb{Z}$ ,  $\omega$ , or  $\omega^*$ , the  $\phi$ -image  $X_a = \phi''[a]_{E_0}$  of the  $E_0$ -class of any  $a \in 2^\omega$  is  $\leq_{\text{lex}}$ -ordered similarly to a subset of  $\mathbb{Z}$ . If  $X_a = \{x_a\}$  is a singleton then put  $\psi(a) = x_a$ .

Assume now that  $X_a$  contains at least two points. In this case we can effectively pick an element in  $X_a$ ! Indeed, there is a maximal sequence  $u \in 2^{<\alpha}$  such that  $u \subseteq x$  for each  $x \in X_a$ . Then the set  $X_a^{\text{left}} = \{x \in X : u \wedge 0 \subseteq x\}$  contains a  $\leq_{\text{lex}}$ -largest element, which we denote by  $\psi(a)$ .

To conclude,  $\psi$  is a Borel reduction of  $E_0$  to the equality on  $2^\alpha$ , i.e.  $a E_0 b$  iff  $\psi(a) = \psi(b)$ , which is impossible because  $E_0$  is not a smooth Borel equivalence relation (see [1]).

**2. The dichotomy.** As usual, it will be assumed that the p.q.-o.  $\preceq$  of Theorem 3 is a  $\Delta_1^1$  relation. Let  $\approx$  denote the associated equivalence.

Following [2] let, for  $\alpha < \omega_1^{\text{CK}}$ ,  $\mathcal{F}_\alpha$  be the family of all h.o.p.  $\Delta_1^1$  functions  $f : \langle \mathcal{N}; \preceq \rangle \rightarrow \langle 2^\alpha; \leq_{\text{lex}} \rangle$ . Then  $\mathcal{F} = \bigcup_{\alpha < \omega_1^{\text{CK}}} \mathcal{F}_\alpha$  is a (countable)  $\Pi_1^1$  set, in a suitable coding system for functions of this type. (See [2] for details.)

Define, for  $x, y \in \mathcal{N}$ ,  $x \equiv y$  iff  $f(x) = f(y)$  for any  $f \in \mathcal{F}$ .

LEMMA 4 (see [2]).  $\equiv$  is a  $\Sigma_1^1$  equivalence relation including  $\approx$ .

Proof. As  $\preceq$  is  $\Delta_1^1$ , one gets by a rather standard argument a  $\Pi_1^1$  set  $N \subseteq \omega$  and a function  $f_n \in \mathcal{F}$  for any  $n \in N$  so that  $\mathcal{F} = \{f_n : n \in N\}$  and the relations  $n \in N \wedge f_n(x) \leq_{\text{lex}} f_n(y)$  and  $n \in N \wedge f_n(x) <_{\text{lex}} f_n(y)$  are presentable in the form  $n \in N \wedge \mathcal{O}(x, y)$  and  $n \in N \wedge \mathcal{O}'(x, y)$  where  $\mathcal{O}, \mathcal{O}'$  are  $\Sigma_1^1$  relations. Now  $x \equiv y$  iff  $\forall n (n \in N \Rightarrow f_n(x) = f_n(y))$ , as required. ■

CASE 1:  $\equiv$  coincides with  $\approx$ . Let us show how this implies (I) of Theorem 3. The set

$$P = \{\langle x, y, n \rangle : x \not\approx y \wedge f_n(x) \neq f_n(y)\}$$

is  $\Pi_1^1$  and, by the assumption of Case 1, its projection on  $x, y$  coincides with the complement of  $\approx$ . Let  $Q \subseteq P$  be a  $\Pi_1^1$  set uniformizing  $P$  in the sense of  $\mathcal{N}^2 \times \omega$ . Then  $Q$  is  $\Delta_1^1$  because

$$Q(x, y, n) \Leftrightarrow x \not\approx y \wedge \forall n' \neq n (\neg Q(x, y, n')).$$

It follows that  $N' = \{n : \exists x, y Q(x, y, n)\} \subseteq N$  is  $\Sigma_1^1$ . Therefore by the  $\Sigma_1^1$  separation theorem there is a  $\Delta_1^1$  set  $M$  such that  $N' \subseteq M \subseteq N$  <sup>(6)</sup>.

Consider a  $\Delta_1^1$  enumeration  $M = \{n_l : l \in \omega\}$ . For any  $l$ ,  $f_{n_l} \in \mathcal{F}_\alpha$  for some ordinal  $\alpha = \alpha_l < \omega_1^{\text{CK}}$ . Another standard argument (see

<sup>(6)</sup> Harrington *et al.* [2] use a general reflection theorem to get such a set, but a more elementary reasoning sometimes has advantage.

[2]) shows that in this case (e.g. when  $M \subseteq N$  is a  $\Delta_1^1$  set) the ordinals  $\alpha_l$  are bounded by some  $\alpha < \omega_1^{\text{CK}}$ . It follows that the function  $h(x) = f_{n_0}(x) \wedge f_{n_1}(x) \wedge f_{n_2}(x) \wedge \dots \wedge f_{n_l}(x) \wedge \dots$  belongs to some  $\mathcal{F}_\beta$ ,  $\beta \leq \alpha \cdot \omega$ . On the other hand, by the construction we have  $x \approx y \Leftrightarrow h(x) = h(y)$ , hence  $h$  satisfies (I) of Theorem 3.

CASE 2:  $\approx \not\subseteq \equiv$ . Assuming this we work towards (II) of Theorem 3.

**3. The domain of singularity.** By the assumption the  $\Sigma_1^1$  set  $A = \{x : \exists y (x \approx y \wedge x \not\equiv y)\}$  is non-empty.

Define  $X \equiv Y$  iff we have  $\forall x \in X \exists y \in Y (x \equiv y)$  and *vice versa*.

PROPOSITION 5. *Let  $X, Y \subseteq A$  be non-empty  $\Sigma_1^1$  sets satisfying  $X \equiv Y$ . Then the sets*

$$P_+ = \{\langle x, y \rangle \in X \times Y : x \equiv y \wedge x \preceq y\}, \quad \text{and}$$

$$P_- = \{\langle x, y \rangle \in X \times Y : x \equiv y \wedge x \not\preceq y\}$$

are non-empty  $\Sigma_1^1$  sets, their projections  $(^7)$   $\text{pr}_1 P^+$  and  $\text{pr}_1 P^-$  are  $\Sigma_1^1$ -dense in  $X$   $(^8)$ , while the projections  $\text{pr}_2 P^+$  and  $\text{pr}_2 P^-$  are  $\Sigma_1^1$ -dense in  $Y$ .

PROOF. The density easily follows from the non-emptiness, so let us concentrate on the latter. We prove that  $P_+ \neq \emptyset$ .

Suppose on the contrary that  $P_+ = \emptyset$ . Then there is a single function  $f \in \mathcal{F}$  such that the set  $\{\langle x, y \rangle \in X \times Y : f(x) = f(y) \wedge x \preceq y\}$  is empty. (See the reasoning in Case 1 of Section 2.) Define

$$X_\infty = \{x : \forall y \in Y (f(x) = f(y) \Rightarrow x \not\preceq y)\},$$

so that  $X_\infty$  is a  $\Pi_1^1$  set and  $X \subseteq X_\infty$  but  $Y \cap X_\infty = \emptyset$ . Using separation, we can easily define an increasing sequence of sets

$$X = X_0 \subseteq U_0 \subseteq X_1 \subseteq U_1 \subseteq \dots \subseteq X_n \subseteq U_n \subseteq \dots \subseteq X_\infty$$

so that  $U_n = \{x' : \exists x \in X_n (f(x) = f(x') \wedge x \preceq x')\}$  while  $X_{n+1} \in \Delta_1^1$  for all  $n$ . (Note that if  $X_n \subseteq X_\infty$  and  $U_n$  is defined as indicated then  $U_n \subseteq X_\infty$  too.) Moreover, a proper execution of the construction  $(^9)$  allows getting the final set  $U = \bigcup_n U_n = \bigcup_n X_n$  in  $\Delta_1^1$ . Note that  $X \subseteq U$ , but  $Y \cap U = \emptyset$  since  $U \subseteq X_\infty$ .

Put  $f'(x) = f(x) \wedge 1$  whenever  $x \in U$ , and  $f'(x) = f(x) \wedge 0$  otherwise. We assert that  $f' \in \mathcal{F}$ . Indeed, suppose that  $x' \preceq y'$ ; we prove  $f'(x') \leq_{\text{lex}} f'(y')$ .

$(^7)$  For a set  $P \subseteq \mathcal{N}^2$ ,  $\text{pr}_1 P$  and  $\text{pr}_2 P$  have the obvious meaning of the projections on the resp. 1st and 2nd copy of  $\mathcal{N}$ .

$(^8)$  That is, each of them intersects any non-empty  $\Sigma_1^1$  set  $X' \subseteq X$ .

$(^9)$  We refer to the proof of an “invariant” effective separation theorem in [1], which includes a similar construction.

It can be assumed that  $f(x') = f(y')$ . It remains to check that  $x' \in U \Rightarrow y' \in U$ , which easily follows from the definition of the sets  $U_n$ . Thus  $f' \in \mathcal{F}$ .

However, clearly  $f'(x) \neq f'(y)$ , hence  $x \not\equiv y$ , whenever  $x \in X$  and  $y \in Y$ , which contradicts the assumption that  $X \equiv Y$ .

Now we prove that  $P_- \neq \emptyset$ . Consider first the case  $X = Y$ . Suppose on the contrary that  $P_- = \emptyset$ . Then, as above, there is a single function  $f \in \mathcal{F}$  such that the set  $\{\langle x, y \rangle \in X^2 : f(x) = f(y) \wedge x \not\preceq y\}$  is empty, so that  $\equiv$  and  $\approx$  coincide on  $X$ . Our plan is to find functions  $f', f'' \in \mathcal{F}$  such that

$$Q' = \{\langle x, y \rangle \in X \times \mathcal{N} : f'(x) = f'(y) \wedge y \not\preceq x\},$$

$$Q'' = \{\langle x, y \rangle \in X \times \mathcal{N} : f''(x) = f''(y) \wedge x \not\preceq y\}$$

are empty sets; then  $Q = \{\langle x, y \rangle \in X \times \mathcal{N} : x \equiv y \wedge y \not\preceq x\} = \emptyset$ , which contradicts  $\emptyset \neq X \subseteq A$ .

Let us find  $f'$ ; the case of the other function is similar. Define

$$X_\infty = \{x : \forall x' \in X (f(x) = f(x') \Rightarrow x \preceq x')\},$$

so that  $X_\infty$  is  $\Pi_1^1$  and  $X \subseteq X_\infty$ . As above there is a sequence of sets

$$X = X_0 \subseteq U_0 \subseteq X_1 \subseteq U_1 \subseteq \dots \subseteq X_n \subseteq U_n \subseteq \dots \subseteq X_\infty$$

such that  $U_n = \{u : \exists x \in X_n (f(x) = f(u) \wedge u \preceq x)\}$  while  $X_{n+1} \in \Delta_1^1$  for all  $n$  and the final set  $U = \bigcup_n U_n = \bigcup_n X_n$  belongs to  $\Delta_1^1$ .

Set  $f'(x) = f(x) \wedge 0$  whenever  $x \in U$ , and  $f'(x) = f(x) \wedge 1$  otherwise. Then  $f' \in \mathcal{F}$ . We prove that  $f'$  witnesses that  $Q' = \emptyset$ . Consider any  $x \in X$  and  $y \in \mathcal{N}$  such that  $f'(x) = f'(y)$ . Then in particular  $f(x) = f(y)$  and  $x \in U \Leftrightarrow y \in U$ , so that  $y \in U$  because we know that  $x \in X \subseteq U$ . Thus  $y \in X_\infty$ , so by definition  $y \preceq x$ , as required.

Finally, we prove  $P_- \neq \emptyset$  in the general case. By the result for the case  $X = Y$ , the  $\Sigma_1^1$  set  $P' = \{\langle x, x' \rangle \in X^2 : x \equiv x' \wedge x \not\preceq x'\}$  is non-empty. Let  $X' = \{x' \in X : \exists x P'(x, x')\}$  and  $Y' = \{y \in Y : \exists x' \in X' (x' \equiv y)\}$ , so that  $X', Y'$  are  $\Sigma_1^1$  sets satisfying  $X' \equiv Y'$ . By the result for  $P_+$  there exist  $x' \in X'$  and  $y \in Y'$  satisfying  $x' \equiv y$  and  $y \preceq x'$ . Now there is  $x \in X$  such that  $x \equiv x'$  and  $x \not\preceq x'$ . Then  $x \equiv y$  and  $x \not\preceq y$ , as required. ■

**4. The forcing notions involved.** Our further strategy will be the following. We shall define a generic extension of the universe  $\mathbf{V}$  (where Theorem 3 is being proved) in which there exists a function  $F$  which witnesses (II) of Theorem 3. However, as the existence of such a function is a  $\Sigma_2^1$  statement, we obtain the result for  $\mathbf{V}$  by the Shoenfield absoluteness theorem <sup>(10)</sup>.

DEFINITION 6.  $\mathbb{P}$  is the collection of all non-empty  $\Sigma_1^1$  sets  $X \subseteq A$ .

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<sup>(10)</sup> In fact, the proof can be conducted without any use of metamathematics, as in [1], but at the cost of longer reasoning.

It is a standard fact that  $\mathbb{P}$  (the *Gandy forcing*) forces a real which is the only real which belongs to every set in the generic set  $G \subseteq \mathbb{P}$ . (We identify  $\Sigma_1^1$  sets in the ground universe  $\mathbf{V}$  with their copies in the extension.)

DEFINITION 7.  $\mathbb{P}_2^+$  is the collection of all non-empty  $\Sigma_1^1$  sets  $P \subseteq A^2$  such that  $P(x, y) \Rightarrow x \equiv y \wedge x \preceq y$ . The collection  $\mathbb{P}_2^-$  is defined similarly but with the requirement  $P(x, y) \Rightarrow x \equiv y \wedge x \not\preceq y$  instead.

Both  $\mathbb{P}_2^+$  and  $\mathbb{P}_2^-$  are non-empty forcing notions by Proposition 5. Each of them forces a pair of reals  $\langle x, y \rangle \in A^2$  satisfying resp.  $x \preceq y$  and  $x \not\preceq y$ .

DEFINITION 8.  $\mathbb{P}_{\equiv}^2$  is the collection of all sets of the form  $\mathcal{T} = X \times Y$  where  $X, Y$  are sets in  $\mathbb{P}$  satisfying  $X \equiv Y$ .

LEMMA 9.  $\mathbb{P}_{\equiv}^2$  forces a pair of reals  $\langle x, y \rangle$  such that  $x \not\preceq y$ .

Proof. Suppose that, on the contrary, a condition  $\mathcal{T}_0 = X_0 \times Y_0$  in  $\mathbb{P}_{\equiv}^2$  forces  $x \preceq y$ . Consider a more complicated forcing  $\mathfrak{P}$  which consists of forcing conditions of the form  $\mathfrak{p} = \langle \mathcal{T}, P, \mathcal{T}', Q \rangle$ , where  $\mathcal{T} = X \times Y$  and  $\mathcal{T}' = X' \times Y'$  belong to  $\mathbb{P}_{\equiv}^2$ ,  $P \in \mathbb{P}_2^+$ ,  $P \subseteq Y \times X'$ ,  $Q \in \mathbb{P}_2^-$ ,  $Q \subseteq X \times Y'$ , and the sets  $\text{pr}_1 P \subseteq Y$ ,  $\text{pr}_2 P \subseteq X'$ ,  $\text{pr}_1 Q \subseteq X$  and  $\text{pr}_2 Q \subseteq Y'$  are  $\Sigma_1^1$ -dense in resp.  $Y, X', X, Y'$ .

For instance, setting  $P_0 = \{\langle y, x' \rangle \in Y_0 \times X_0 : y \equiv x' \wedge y \preceq x'\}$  and  $Q_0 = \{\langle x, y' \rangle \in X_0 \times Y_0 : x \equiv y' \wedge x \not\preceq y'\}$ , we get a condition  $\mathfrak{p}_0 = \langle \mathcal{T}_0, P_0, \mathcal{T}_0, Q_0 \rangle \in \mathfrak{P}$  by Proposition 5.

It is the principal fact that if  $\mathfrak{p} = \langle \mathcal{T}, P, \mathcal{T}', Q \rangle \in \mathfrak{P}$  and we strengthen one of the components within the corresponding forcing notion then this can be appropriately reflected in the other components. To be concrete assume that, for instance,  $P^* \in \mathbb{P}_2^+$ ,  $P^* \subseteq P$ , and find a condition  $\mathfrak{p}_1 = \langle \mathcal{T}_1, P_1, \mathcal{T}'_1, Q_1 \rangle \in \mathfrak{P}$  satisfying  $\mathcal{T}_1 \subseteq \mathcal{T}$ ,  $\mathcal{T}'_1 \subseteq \mathcal{T}'$ ,  $P_1 \subseteq P^*$ , and  $Q_1 \subseteq Q$ .

Assume that  $\mathcal{T} = X \times Y$  and  $\mathcal{T}' = X' \times Y'$ . Consider the non-empty  $\Sigma_1^1$  sets  $Y_2 = \text{pr}_1 P^* \subseteq Y$  and  $X_2 = \{x \in X : \exists y \in Y_2 (x \equiv y)\}$ . It follows from Proposition 5 that  $Q_1 = \{\langle x, y \rangle \in Q : x \in X_2\} \neq \emptyset$ , hence  $Q_1$  is a condition in  $\mathbb{P}_2^-$  and  $X_1 = \text{pr}_1 Q_1$  is a non-empty  $\Sigma_1^1$  subset of  $X_2 \subseteq X$ .

The set  $Y_1 = \{y \in Y_2 : \exists x \in X_1 (x \equiv y)\}$  satisfies  $X_1 \equiv Y_1$ , therefore  $\mathcal{T}_1 = X_1 \times Y_1 \in \mathbb{P}_{\equiv}^2$ . Furthermore,  $P_1 = \{\langle y, x \rangle \in P^* : y \in Y_1\} \in \mathbb{P}_2^+$ .

Put  $X'_1 = \text{pr}_2 P_1 \subseteq X'$  and  $Y'_1 = \text{pr}_2 Q_1 \subseteq Y'$ . Notice that  $Y_1 \equiv X'_1$  because any condition in  $\mathbb{P}_2^+$  is a subset of  $\equiv$ , similarly  $X_1 \equiv Y'_1$ , and  $X_1 \equiv Y_1$  (see above). It follows that  $X'_1 \equiv Y'_1$ , hence  $\mathcal{T}'_1 = X'_1 \times Y'_1$  is a condition in  $\mathbb{P}_{\equiv}^2$ .

Now  $\mathfrak{p}_1 = \langle \mathcal{T}_1, P_1, \mathcal{T}'_1, Q_1 \rangle \in \mathfrak{P}$  as required.

We conclude that  $\mathfrak{P}$  forces “quadruples” of reals  $\langle x, y, x', y' \rangle$  such that the pairs  $\langle x, y \rangle$  and  $\langle x', y' \rangle$  are  $\mathbb{P}_{\equiv}^2$ -generic, hence satisfy  $x \preceq y$  and  $x' \preceq y'$  provided the generic set contains  $\mathcal{T}_0$ —by the assumption above. Furthermore, the pair  $\langle y, x' \rangle$  is  $\mathbb{P}_2^+$ -generic, hence  $y \preceq x'$ , while the pair  $\langle x, y' \rangle$  is  $\mathbb{P}_2^-$ -generic, hence  $x \not\preceq y'$ , which is a contradiction. ■

**5. The splitting construction.** Let, in the universe  $\mathbf{V}$ ,  $\kappa = 2^{\aleph_0}$ . Let  $\mathbf{V}^+$  be a  $\kappa$ -collapse extension of  $\mathbf{V}$ .

Our aim is to define, in  $\mathbf{V}^+$ , a splitting system of sets which leads to a function  $F$  satisfying (II) of Theorem 3. Let us fix two points before the construction starts.

*First*, as the forcing notions involved are countable in  $\mathbf{V}$ , there exist, in  $\mathbf{V}^+$ , enumerations  $\{D(n) : n \in \omega\}$ ,  $\{D_2(n) : n \in \omega\}$ , and  $\{D^2(n) : n \in \omega\}$  of all open dense sets in resp.  $\mathbb{P}$ ,  $\mathbb{P}_2^+$ ,  $\mathbb{P}_{\equiv}^2$ , which (the dense sets) belong to  $\mathbf{V}$ , such that  $D(n+1) \subseteq D(n)$  etc. for each  $n$ .

*Second*, we introduce the notion of a crucial pair. A pair  $\langle u, v \rangle$  of binary sequences  $u, v \in 2^n$  is called *crucial* iff  $u = 1^k \wedge 0 \wedge w$  and  $v = 0^k \wedge 1 \wedge w$  for some  $k < n$  and  $w \in 2^{n-k-1}$ . One easily sees that the graph of all crucial pairs in  $2^n$  is actually a chain connecting all members of  $2^n$ .

We define, in  $\mathbf{V}^+$ , a system of sets  $X_u \in \mathbb{P}$ , where  $u \in 2^{<\omega}$ , and sets  $P_{uv} \in \mathbb{P}_2^+$ ,  $\langle u, v \rangle$  being a crucial pair in some  $2^n$ , satisfying the following conditions:

- (1)  $X_u \in D(n)$  whenever  $u \in 2^n$ ;  $X_{u \wedge i} \subseteq X_u$ ;
- (2) if  $\langle u, v \rangle$  is a crucial pair in  $2^n$  then  $P_{uv} \in D_2(n)$  and  $P_{u \wedge i, v \wedge i} \subseteq P_{uv}$ ;
- (3) if  $u, v \in 2^n$  and  $u(n-1) \neq v(n-1)$  then  $X_u \times X_v \in \mathbb{P}_{\equiv}^2$ ,  $X_u \times X_v \in D^2(n)$ , and  $X_u \cap X_v = \emptyset$ ;
- (4) if  $\langle u, v \rangle$  is a crucial pair in  $2^n$  then  $\text{pr}_1 P_{uv} = X_u$  and  $\text{pr}_2 P_{uv} = X_v$ .

*Why does this imply the existence of a required function?* First of all for any  $a \in 2^\omega$  (in  $\mathbf{V}^+$ ) the sequence of sets  $X_{a \upharpoonright n}$  is  $\mathbb{P}$ -generic over  $\mathbf{V}$  by (1), therefore the intersection  $\bigcap_{n \in \omega} X_{a \upharpoonright n}$  is a singleton. Let  $F(a) \in \mathcal{N}$  be its only element.

It does not take much effort to prove that  $F$  is continuous and 1-1.

Consider  $a, b \in 2^\omega$  satisfying  $a \not\leq_0 b$ . Then  $a(n) \neq b(n)$  for infinitely many  $n$ , hence the pair  $\langle F(a), F(b) \rangle$  is  $\mathbb{P}_{\equiv}^2$ -generic by (3), thus  $F(a)$  and  $F(b)$  are  $\preceq$ -incomparable by Lemma 9.

Consider  $a, b \in 2^\omega$  satisfying  $a \leq_0 b$ . We may assume that  $a$  and  $b$  are  $\leq_0$ -neighbours, i.e.  $a = 1^k \wedge 0 \wedge c$  while  $b = 0^k \wedge 1 \wedge c$  for some  $k \in \omega$  and  $c \in 2^\omega$ . Then by (2) the sequence of sets  $P_{a \upharpoonright n, b \upharpoonright n}$ ,  $n > k$ , is  $\mathbb{P}_2^+$ -generic, hence it results in a pair of reals satisfying  $x \preceq y$ . However,  $x = F(a)$  and  $y = F(b)$  by (4).

*The construction of a splitting system. We argue in  $\mathbf{V}^+$ .*

Suppose that the construction has been completed up to a level  $n$ ; we will expand it to the next level. From now on  $s, t$  will denote sequences in  $2^n$  while  $u, v$  will denote sequences in  $2^{n+1}$ .

To start with, we set  $X_{s \wedge i} = X_s$  for all  $s \in 2^n$  and  $i = 0, 1$ , and  $P_{s \wedge i, t \wedge i} = P_{st}$  whenever  $i = 0, 1$  and  $\langle s, t \rangle$  is a crucial pair in  $2^n$ .

For the “initial” crucial pair  $\langle 1^{n\wedge 0}, 0^{n\wedge 1} \rangle$  at this level let  $P_{1^{n\wedge 0}, 0^{n\wedge 1}} = X_{1^{n\wedge 0}} \times X_{0^{n\wedge 1}} = X_{1^n} \times X_{0^n}$ . Then  $P_{1^{n\wedge 0}, 0^{n\wedge 1}} \in \mathbb{P}_{\equiv}^2$  <sup>(11)</sup>.

This ends the definition of “initial values” at the  $(n + 1)$ th level. The plan is to gradually “shrink” the sets in order to fulfill the requirements.

STEP 1. We take care of item (1). Consider an arbitrary  $u_0 = s_0^{\wedge i} \in 2^{n+1}$ . As  $D(n)$  is dense there is a set  $X' \in D(n)$  with  $X' \subseteq X_{u_0}$ . The intention is to take  $X'$  as the “new”  $X_{u_0}$ . But this change has to be expanded through the chain of crucial pairs, in order to preserve (4).

Thus put  $X'_{u_0} = X'$ . Suppose that  $X'_u$  has been defined and is included in  $X_u$ , the “old” version, for some  $u \in 2^{n+1}$ , and  $\langle u, v \rangle$  is a crucial pair,  $v \in 2^{n+1}$  being not yet encountered. Define  $P'_{uv} = (X'_u \times \mathcal{N}) \cap P_{uv}$  and  $X'_v = \text{pr}_2 P'_{uv}$ . Clearly (4) holds for the “new” sets  $X'_u$ ,  $X'_v$ , and  $P'_{uv}$ .

The construction describes how the original change from  $X_{u_0}$  to  $X'_{u_0}$  spreads through the chain of crucial pairs in  $2^{n+1}$ , resulting in a system of new sets,  $X'_u$  and  $P'_{uv}$ , which satisfy (1) for the particular  $u_0 \in 2^{n+1}$ . We iterate this construction consecutively for all  $u_0 \in 2^{n+1}$ , getting finally a system of sets satisfying (1) (fully) and (4), which we shall denote by  $X_u$  and  $P_{uv}$  from now on.

STEP 2. We take care of item (3). Fix a pair of  $u_0$  and  $v_0$  in  $2^{n+1}$  such that  $u_0(n) = 0$  and  $v_0(n) = 1$ . By the density of  $D^2(n)$ , there is a set  $X'_{u_0} \times X'_{v_0} \in D^2(n)$  included in  $X_{u_0} \times X_{v_0}$ . We may assume that  $X'_{u_0} \cap X'_{v_0} = \emptyset$ . (Indeed, it easily follows from Proposition 5, for  $P_-$ , that there exist reals  $x_0 \in X_{u_0}$  and  $y_0 \in X_{v_0}$  satisfying  $x_0 \equiv y_0$  but  $x_0 \neq y_0$ , say  $x_0(k) = 0$  while  $y_0(k) = 1$ . Define

$$X = \{x \in X_0 : x(k) = 0 \wedge \exists y \in Y_0 (y(k) = 1 \wedge x \equiv y)\},$$

and  $Y$  correspondingly; then  $X \equiv Y$  and  $X \cap Y = \emptyset$ .)

Spread the change from  $X_{u_0}$  to  $X'_{u_0}$  and from  $X_{v_0}$  to  $X'_{v_0}$  through the chain of crucial pairs in  $2^{n+1}$ , by the method of Step 1, until the wave of spreading from  $u_0$  meets the wave of spreading from  $v_0$  at the “meeting” crucial pair  $\langle 1^{n\wedge 0}, 0^{n\wedge 1} \rangle$ . This leads to a system of sets  $X'_u$  and  $P'_{uv}$  which satisfy (3) for the particular pair  $\langle u_0, v_0 \rangle$  and still satisfy (4) possibly except for the “meeting” crucial pair  $\langle 1^{n\wedge 0}, 0^{n\wedge 1} \rangle$  (for which basically the set  $P'_{1^{n\wedge 0}, 0^{n\wedge 1}}$  is not yet defined at this step).

Note that Step 1 leaves  $P_{1^{n\wedge 0}, 0^{n\wedge 1}}$  in the form  $X_{1^{n\wedge 0}} \times X_{0^{n\wedge 1}}$  (where  $X_{1^{n\wedge 0}}$  and  $X_{0^{n\wedge 1}}$  are the “versions” at the end of Step 1). We now have the “new” sets,  $X'_{1^{n\wedge 0}}$  and  $X'_{0^{n\wedge 1}}$ , included in resp.  $X_{1^{n\wedge 0}}$  and  $X_{0^{n\wedge 1}}$  and satisfying  $X'_{0^{n\wedge 1}} \equiv X'_{1^{n\wedge 0}}$  (because we had  $X'_{u_0} \equiv X'_{v_0}$  at the beginning of

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<sup>(11)</sup> It easily follows from (2) and (4) that  $X_s \equiv X_t$  for all  $s, t \in 2^n$ , because  $s$  and  $t$  are connected in  $2^n$  by a unique chain of crucial pairs.

the change). It remains to define  $P'_{1^{n \wedge 0}, 0^{n \wedge 1}} = X'_{1^{n \wedge 0}} \times X'_{0^{n \wedge 1}}$ . This ends the consideration of the pair  $\langle u_0, v_0 \rangle$ .

Applying this construction consecutively for all pairs of  $u_0 \in P_0$  and  $v_0 \in P_1$  (including the pair  $\langle 1^{n \wedge 0}, 0^{n \wedge 1} \rangle$ ) we finally get a system of sets satisfying (1), (3), and (4), which will be denoted still by  $X_u$  and  $P_{uv}$ .

STEP 3. We finally take care of (2). Consider a crucial pair  $\langle u_0, v_0 \rangle$  in  $2^{n+1}$ . By density, there exists a set  $P'_{u_0, v_0} \in D_2(n)$  with  $P'_{u_0, v_0} \subseteq P_{u_0, v_0}$ . (In the case when  $\langle u_0, v_0 \rangle$  is the pair  $\langle 1^{n \wedge 0}, 0^{n \wedge 1} \rangle$  we rather apply Proposition 5 to obtain the set  $P'_{u_0, v_0}$ .)

Define  $X'_{u_0} = \text{pr}_1 P'_{u_0, v_0}$  and  $X'_{v_0} = \text{pr}_2 P'_{u_0, v_0}$  and spread this change through the chain of crucial pairs in  $2^{n+1}$ . (Note that  $X'_{u_0} \equiv X'_{v_0}$  as sets in  $\mathbb{P}^2_{\equiv}$  are included in  $\equiv$ . This keeps  $X'_u \equiv X'_v$  for all  $u, v \in 2^{n+1}$  through the spreading.)

Executing this step for all crucial pairs in  $2^{n+1}$ , we finally end the construction, in  $\mathbf{V}^+$ , of a system of sets satisfying (1) through (4). ■ Theorem 3

### References

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