

**A recursion-theoretic characterization of instances of  $B\Sigma_n$  provable  
in  $\Pi_{n+1}(N)$**

by

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**Abstract.** Let  $B\psi$  denote the sentence

$$(\forall a, t)((\forall x)_{\leq t}(\exists y)\psi(x, a, y) \Rightarrow (\exists z)(\forall x)_{\leq t}(\exists y)_{< z}\psi(x, a, y))$$

where  $\psi$  is  $\Pi_{n-1}$ ,  $n \geq 1$ .

We give a recursion-theoretic characterization of those  $\Pi_{n-1}$  formulas  $\psi$  that  $B\psi$  is provable in  $\Pi_{n+1}(N)$ .

As a corollary we infer that  $\Pi_{n+1}(N)$  does not prove  $B\Sigma_n$ .

The independence of  $B\Sigma_n$  from the theory of true  $\Pi_{n+1}$  sentences —  $\Pi_{n+1}(N)$  — was first proved by Parsons in [P] by a combination of a proof-theoretic and recursion-theoretic argument. There is also a model-theoretic proof possible. Paris and Kirby in [PK] prove that  $B\Sigma_n$  is independent from  $\text{I}\Sigma_{n-1}$  in a model theoretic way. One can modify their argument so as to show the unprovability of  $B\Sigma_n$  in  $\Pi_{n+1}(N)$ .

Note that  $\Pi_{n+1}(N)$  is a stronger theory than  $\text{I}\Sigma_{n-1}$ .

We give another proof of the independence of  $B\Sigma_n$  from  $\Pi_{n+1}(N)$ . It is similar to Parsons's proof, but in addition to the independence result, it gives a full characterization of formulas  $\psi$  such that the corresponding instance of  $B\Sigma_n$  is provable in  $\Pi_{n+1}(N)$ . To formulate our characterization theorem we need two definitions of recursion-theoretic character.

**DEFINITION 1.** A relation  $R(a, x) \subseteq N \times N$  is called *almost a function* if there is a number  $L \in N$  such that for every  $a$ , the set  $\{x: R(a, x)\}$  has at most  $L$  elements.

**DEFINITION 2.** A relation  $\hat{R}(a, x)$  *uniformizes*  $R(a, x)$  if for every  $a$

$$(\exists x) R(a, x) \Rightarrow (\exists x) (R(a, x) \ \& \ \hat{R}(a, x)).$$

Our main theorem is the following:

**THEOREM 1.**  $\Pi_{n+1}(N) \vdash B\psi$  where  $\psi$  is  $\Pi_{n-1}$  iff the  $\Pi_n$  relation  $R(\langle t, a \rangle, x)$ :

$$(\forall y) \neg \psi(x, a, y) \ \& \ x \leq t$$

is uniformizable in  $N$  by a  $\Sigma_n$  almost function.

Proof. First we prove the theorem in the " $\Leftarrow$ " direction, which is easier. Let  $\psi$  always denote a  $\Pi_{n-1}$  formula.

LEMMA 1. Assume that there is a  $\Sigma_{n+1}$  formula  $\varphi(t, a)$  such that

$$\Pi_{n+1}(N) \vdash (\forall t, a)((\forall x)_{\leq t}(\exists y)\psi(x, a, y) \Leftrightarrow \varphi(t, a)).$$

Then  $\Pi_{n+1}(N) \vdash B\psi$ .

Proof. By the assumption,  $B\psi$  is equivalent in  $\Pi_{n+1}(N)$  to

$$(\forall t, a)(\varphi(t, a) \Rightarrow (\exists z)(\forall x)_{\leq t}(\exists y)_{\leq z}\psi(x, a, y)).$$

The above sentence is  $\Pi_{n+1}$  and true, hence it is provable in  $\Pi_{n+1}(N)$ . Thus  $\Pi_{n+1}(N) \vdash B\psi$ . ■

LEMMA 2. Assume that  $R(a, x)$  is  $\Pi_n$  and is uniformizable by a  $\Sigma_n$  almost function. Then the formula  $\varphi(a): (\exists x)R(a, x)$  is provably  $\Pi_{n+1}$  in  $\Pi_{n+1}(N)$ , i.e. there is a  $\Pi_{n+1}$  formula  $\hat{\varphi}(a)$  such that  $\Pi_{n+1}(N) \vdash (\forall a)(\varphi(a) \Leftrightarrow \hat{\varphi}(a))$ .

Proof. Let  $\hat{R}(a, x)$  be a  $\Sigma_n$  almost function uniformizing  $R$ . Let  $L$  be given by Definition 1 for  $\hat{R}$ . We have

$$\begin{aligned} (\forall a)\{(\exists x)R(a, x) \Leftrightarrow \bigvee_{i=1}^L [(\exists x_1 \dots x_i)(\bigwedge_{j,k=1}^i x_j \neq x_k \ \& \ \bigwedge_{j=1}^i \hat{R}(a, x_j)) \ \& \\ \& \ (\forall x_1 \dots x_i)((\bigwedge_{j,k=1}^i x_j \neq x_k \ \& \ \bigwedge_{j=1}^i \hat{R}(a, x_j)) \Rightarrow \bigvee_{j=1}^i R(a, x_j)]]\}. \end{aligned}$$

Then the " $\Rightarrow$ " part of the above equivalence is a  $\Pi_{n+1}$  sentence and, being true, it is provable in  $\Pi_{n+1}(N)$ . The " $\Leftarrow$ " part is provable in logic.

Moreover, the right hand side of the equivalence is a  $\Pi_{n+1}$  formula (in fact it is a boolean combination of  $\Sigma_n$  formulas).

Thus  $\varphi(a)$  is equivalent in  $\Pi_{n+1}(N)$  to the  $\Pi_{n+1}$  formula

$$\begin{aligned} \bigvee_{i=1}^L [(\exists x_1 \dots x_i)(\bigwedge_{j,k=1}^i x_j \neq x_k \ \& \ \bigwedge_{j=1}^i \hat{R}(a, x_j)) \ \& \ (\forall x_1 \dots x_i)((\bigwedge_{j,k=1}^i x_j \neq x_k \ \& \\ \& \ \bigwedge_{j=1}^i \hat{R}(a, x_j)) \Rightarrow \bigvee_{j=1}^i R(a, x_j)]]. \quad \blacksquare \end{aligned}$$

Proof of the " $\Leftarrow$ " part of the theorem. Assume that the relation  $R(\langle t, a \rangle, x): (\forall y)\neg\psi(x, a, y) \ \& \ x \leq t$  is uniformizable by a  $\Sigma_n$  almost function. Then by Lemma 2 the formula  $(\exists x)(\forall y)(\neg\psi(x, a, y) \ \& \ x \leq t)$  is provably  $\Pi_{n+1}$  in  $\Pi_{n+1}(N)$ . Thus, there is a  $\Sigma_{n+1}$  formula  $\varphi(t, a)$  such that  $\Pi_{n+1}(N)$  proves

$$(\forall t, a)((\forall x)_{\leq t}(\exists y)\psi(x, a, y) \Leftrightarrow \varphi(t, a)).$$

Now, by Lemma 1,  $\Pi_{n+1}(N) \vdash B\psi$ .

Proof of the " $\Rightarrow$ " part of the theorem. Consider the following definition.

DEFINITION 3. Let  $\varphi(\vec{x})$  be a formula of the form:

$$(\forall z_0)(\exists w_0) \dots (\forall z_k)(\exists w_k) \varphi'(x, z_0, w_0, \dots, z_k, w_k)$$

where  $\varphi'$  is open.

Since we allow irrelevant quantifiers, an arbitrary formula can be written in this form.

Fix an enumeration of  $Z$ -polynomials. Let  $K \in N$ .

A finite set  $H$  is a  $K$ -closure of  $\vec{t} = \langle t_0 \dots t_m \rangle$  w.r.t.  $\varphi$  if there are sets  $H_0, \dots, H_K$  such that

$$(1) H = H_0 \cup \dots \cup H_K, \quad H_i \subseteq H_{i+1};$$

$$(2) H_0 = \{t_0, \dots, t_m\};$$

(3) if  $i < K$ ,  $r$  is a polynomial of  $l$  variables whose number is less than  $i$  and  $\vec{x} \in H_i$  is of length  $l$  and  $r(\vec{x}) \geq 0$  then  $r(\vec{x}) \in H_{i+1}$  (by  $\vec{x} \in H_i$  we mean that every term of  $\vec{x}$  belongs to  $H_i$ );

(4) there are partial functions  $w_j(z_0 \dots z_j)$  from  $H^{j+1}$  to  $H$  for  $j = 0, \dots, k-1$  such that

(a) if there are  $i_0, \dots, i_j < K$  such that  $i_0 < i_1 < \dots < i_j$  and  $z_0 \in H_{i_0}, \dots, z_j \in H_{i_j}$  then the sentence

$$\begin{aligned} (\forall z_{j+1})(\exists w_{j+1}) \dots (\forall z_k)(\exists w_k) \\ \varphi'(\vec{t}, z_0, w_0(z_0), z_1, w_1(z_0, z_1) \dots z_j, w_j(z_0 \dots z_j), z_{j+1}, w_{j+1} \dots z_k, w_k) \end{aligned}$$

is true.

Remark 1. For every  $K \in N$  there is a number  $L(K) \in N$  such that a minimal  $K$ -closure of any  $\vec{t}$  w.r.t.  $\varphi$  has at most  $L(K)$  elements.

Moreover there is a formula  $\theta_{L(K)}^{\varphi}(\vec{t}, \vec{h})$  of  $L(K) + m$  variables stating the following:  $h_0, \dots, h_{L(K)-1}$  are consecutive, in a certain canonical ordering, elements of a minimal  $K$ -closure of  $\vec{t}$  w.r.t.  $\varphi$ .

We show how to build  $\theta_{L(K)}^{\varphi}$  on a concrete example. Assume that  $\varphi(x)$  is of the form

$$(\forall z_0)(\exists w_0)(\forall z_1)\varphi'(x, z_0, w_0, z_1).$$

Let  $r_0, r_1$  be polynomials of one variable whose numbers are 0,1 respectively. We have:

$$L(0) = 1, \quad \theta_{L(0)}^{\varphi}(t, h_0): h_0 = t,$$

$$L(1) = 3, \quad \theta_{L(1)}^{\varphi}(t, h_0, h_1, h_2): h_0 = t \ \& \ h_1 = r_0(t) \ \& \\ \& \ (\forall z_1)\varphi'(t, h_0, h_2, z_1),$$

$$L(2) = 10, \quad \theta_{L(2)}^{\varphi}(t, h_0 \dots h_9): h_0 = t \ \& \ h_1 = r_0(t) \ \& \ (\forall z_1)\varphi'(t, h_0, h_2, z_1) \\ \& \ h_3 = r_0(h_1) \ \& \ h_4 = r_0(h_2) \ \& \ h_5 = r_1(h_0) \ \& \ h_6 = r_1(h_1) \ \& \\ \& \ h_7 = r_1(h_2) \ \& \ (\forall z_1)\varphi'(t, h_1, h_3, z_1) \ \& \ (\forall z_1)\varphi'(t, h_2, h_9, z_1).$$

Note that, if  $\varphi$  is a  $\Pi_{n+1}$  formula then  $\theta_{L(K)}^{\varphi}$  is  $\Sigma_n$  (even  $\Pi_{n-1}$ ).

Remark 2. We have for any  $K$

$$\text{PA}^- \vdash (\forall t)(\varphi(t) \Rightarrow (\exists \vec{h})\theta_{L(K)}^0(t, \vec{h})).$$

The next lemma is a version of Herbrand's theorem.

LEMMA 3. Let  $\varphi$  be a formula. The following are equivalent:

- (1)  $\text{PA}^- \vdash (\forall \vec{i})\varphi(\vec{i})$ ;
- (2) there is a number  $K \in N$  such that

$$\text{PA}^- \vdash (\forall \vec{i}, \vec{h}) \neg \theta_{L(K)}^0(\vec{i}, \vec{h}).$$

Proof (due to Z. Ratajczyk). Assume (1). Add to the language the constants:  $d_0, \dots, d_m, c_0, c_1, \dots$  and let  $\vec{d} = \langle d_0 \dots d_m \rangle$ ,  $\vec{c}_K = \langle c_0 \dots c_{L(K)-1} \rangle$ . Suppose that (2) does not hold. Then for every  $K$  the theory

$$T_K: \text{PA}^- + \{\theta_{L(0)}^0(\vec{d}, \vec{c}_0), \theta_{L(1)}^0(\vec{d}, \vec{c}_1), \dots, \theta_{L(K)}^0(\vec{d}, \vec{c}_K)\}$$

is consistent.

Then the theory  $T = \bigcup_K T_K$  is consistent, by compactness. Let  $M$  be a model of  $T$  and let  $M' \subseteq M$  be the submodel of  $M$  whose universe consists of interpretations of the constants:

$$t_0, \dots, t_m, h_0, h_1, \dots$$

CLAIM.  $M' \models \text{PA}^- \& \neg \varphi(\vec{i})$  where  $\vec{i} = \langle t_0 \dots t_m \rangle$

Assume that  $\neg \varphi(\vec{i})$  is  $(\forall z_0)(\exists w_0) \dots (\forall z_k)(\exists w_k)\varphi'(\vec{i}, z_0, w_0 \dots z_k, w_k)$ . We have  $M' \models \text{PA}^-$  since it is closed under polynomials in  $M$ . Moreover, for any  $z_0, \dots, z_k \in M$ , there are elements  $w_0(z_0), w_1(z_0, z_1), \dots, w_k(z_0 \dots z_k)$  interpreting appropriate elements of the  $K$ -closures:  $h_0 \dots h_{L(K)-1}$ . Then

$$M' \models \varphi'(\vec{i}, z_0, w_0(z_0), z_1, w_1(z_0, z_1) \dots z_k, w_k(z_0 \dots z_k)).$$

Thus the claim is proved and hence a contradiction. The implication (2)  $\Rightarrow$  (1) is easy. ■

Now we are able to prove the " $\Rightarrow$ " part of the theorem.

Assume that  $\Pi_{n+1}(N)$  proves  $B\psi$ . Especially, there is a  $\Sigma_n$  formula  $\sigma(t, a)$  such that

$$\Pi_{n+1}(N) \vdash (\forall a, t)((\forall x)_{\leq t}(\exists y)\psi(x, a, y) \Leftrightarrow \sigma(t, a)).$$

Let  $\varphi_1, \dots, \varphi_m$  be a finite fragment of  $\Pi_{n+1}(N)$  which proves the above equivalence. Let  $\varphi(a, t)$  be the formula

$$(\varphi_1 \& \dots \& \varphi_m) \Rightarrow ((\forall x)_{\leq t}(\exists y)\psi(x, a, y) \Rightarrow \sigma(t, a)).$$

Then  $\text{PA}^- \vdash (\forall a, t)\varphi(a, t)$ .

By Lemma 3 there is a  $K \in N$  such that

$$\text{PA}^- \vdash (\forall t, a, \vec{h}) \neg \theta_{L(K)}^0(t, a, \vec{h})$$

Here,  $\neg \varphi$  is the formula

$$\varphi_1 \& \dots \& \varphi_m \& (\forall x)_{\leq t}(\exists y)\psi(x, a, y) \& \neg \sigma(t, a).$$

This formula is  $\Pi_{n+1}$ , hence the formulas  $\theta_{L(i)}^0$  are  $\Sigma_n$ . We can assume that  $\theta_{L(i)}^0(t, a, \vec{h})$  define graphs of partial functions of  $t, a$  — otherwise take  $\Sigma_n$  formulas defining graphs of partial functions of  $t, a$  and uniformizing  $\theta_{L(i)}^0$ .

Consider the following relation  $\hat{R}(\langle t, a \rangle, x)$  in  $N \times N$ :

$$\bigwedge_{i=0}^{K-1} (\exists \vec{h}) (\theta_{L(i)}^0(t, a, \vec{h}) \& (x = h_0 \vee x = h_1 \vee \dots \vee x = h_{L(i)-1})).$$

Then  $\hat{R}$  is a  $\Sigma_n$  almost function (to satisfy Definition 1 one can take  $L = \sum_{i < K} L(i)$ ).

CLAIM.  $\hat{R}$  uniformizes  $R$  where  $R(\langle t, a \rangle, x)$  is the relation defined in  $N$  by

$$(\forall y) \neg \psi(x, a, y) \& x \leq t.$$

Proof of the claim. Assume  $(\exists x)R(\langle t, a \rangle, x)$  for a pair  $\langle t, a \rangle \in N$ . Then  $(\exists x)(x \leq t \& (\forall y) \neg \psi(x, a, y))$ . We have  $\neg \sigma(t, a)$  since

$$\sigma(t, a) \Rightarrow (\forall x)_{\leq t}(\exists y)\psi(x, a, y)$$

in  $N$ . Let  $i < K$  be maximal such that  $(\exists \vec{h})\theta_{L(i)}^0(t, a, \vec{h})$ . Such an  $i$  exists since  $(\exists \vec{h})\theta_{L(0)}^0(t, a, \vec{h})$  and  $(\forall \vec{h}) \neg \theta_{L(K)}^0(t, a, \vec{h})$ . Take  $\vec{h}$  such that  $\theta_{L(i)}^0(t, a, \vec{h})$ .

If for every  $h_j \leq t$  there was a  $y$  such that  $\psi(h_j, a, y)$  then, since  $\varphi_1, \dots, \varphi_m$  are true and  $\neg \sigma(t, a)$  is true, we would be able to extend  $\vec{h}$  to an  $(i+1)$ th closure of  $t$ , a w.r.t.  $\neg \varphi$  contradicting the choice of  $i$ . Hence

$$\bigwedge_{j=0}^{L(i)-1} (h_j \leq t \& (\forall y) \neg \psi(h_j, a, y)).$$

The claim has thus been proved and part " $\Rightarrow$ " of the theorem follows. ■

COROLLARY.  $\Pi_{n+1}(N)$  does not prove  $B\Sigma_n$ .

Proof. We shall show that  $\Pi_{n+1}(N)$  does not prove the following form of  $B\Sigma_n$ :

$$(\forall t)((\forall x)_{\leq t}(\exists y)\psi(x, t, y) \Rightarrow (\exists z)(\forall x)_{\leq t}(\exists y)_{< z}\psi(x, t, y))$$

where  $\psi$  runs over  $\Pi_{n-1}$  formulas.

The corollary follows from the following remark:

Remark 5. There is a  $\Pi_n$  relation  $R(t, x)$  in  $N \times N$  such that the relation  $R'(t, x)$  defined as  $R(t, x) \& x \leq t$  is not uniformizable by a  $\Sigma_n$  almost function.

Proof of the remark. Let  $S(s, t, x) \subseteq N \times N \times N$  be a universal  $\Sigma_n$  relation. Let  $R(t, x)$  be  $\neg S(t, t, x)$ . Then  $R'(t, x): R(t, x) \& x \leq t$  is not uniformizable by any  $\Sigma_n$  almost function.

Suppose the converse. Let  $\hat{R}(t, x)$  be a  $\Sigma_n$  almost function and uniformize  $R'$ . Let  $L$  be the bound on the number of  $x$ 's satisfying  $\hat{R}$  with a given  $t$ . Let  $s$  be such that  $s > L$  and

$$(\forall x, t)(\hat{R}(t, x) \Leftrightarrow S(s, t, x)).$$

Let us show that  $(\exists x) R'(s, x)$ . Indeed if  $(\forall x) \neg R'(s, x)$  then  $(\forall x)_{\leq s} \neg R(s, x)$  hence  $(\forall x)_{\leq s} S(s, s, x)$  and hence  $(\forall x)_{\leq s} \hat{R}(s, x)$ . But this contradicts the choice of  $L$ .

Thus  $(\exists x) R'(s, x)$ . Hence, by uniformization  $(\exists x)(R'(s, x) \ \& \ \hat{R}(s, x))$ .

But  $R'(s, x) \Leftrightarrow x \leq s \ \& \ \neg S(s, s, x) \Rightarrow \neg \hat{R}(s, x)$ . Contradiction. Since  $R$  is  $\Pi_n$  it is as required. ■

Proof of the corollary. Let  $\psi$  be a  $\Pi_{n-1}$  formula such that

$$R(t, x) \Leftrightarrow (\forall y) \neg \psi(x, t, y)$$

where  $R$  is taken from the remark. If

$$(\forall t)((\forall x)_{\leq t}(\exists y)\psi(x, t, y) \Rightarrow (\exists z)(\forall x)_{\leq t}(\exists y)_{< z}\psi(x, t, y))$$

was provable in  $\Pi_{n+1}(N)$  then, from the proof of Theorem 1 it follows that  $R(t, x)$  would be uniformizable by a  $\Sigma_n$  almost function and it is not the case. ■

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- [P] C. Parsons, On a number theoretic choice schema and its relation to induction, in *Intuitionisms and Proof Theory*, North Holland 1970, ed. Kino, Myhill, Vesley.

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#### ERRATA

Page line	For	Read
166 <sup>8</sup>	1976	1986
234 <sup>2</sup>	$\text{Pa}^- \vdash (\forall t)(\varphi(t) \Rightarrow (\exists \vec{h}) \Theta_{L(K)}^p(t, \vec{h}))$	$\text{Pa}^- \vdash (\forall \vec{i})(\varphi(\vec{i}) \Rightarrow (\exists \vec{h}) \Theta_{L(K)}^p(t, \vec{h}))$
234 <sup>17</sup>	$M' \Vdash$	$M' \Vdash$
234 <sup>19</sup>	$M' \Vdash$	$M' \Vdash$

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