

## UNIFORMLY ENCLOSING DISCRETIZATION METHODS FOR SEMILINEAR BOUNDARY VALUE PROBLEMS

H.-G. ROOS

*Section of Mathematics, Technical University of Dresden,  
 Dresden, German Democratic Republic*

### 1. Introduction

In the numerical solution of differential equations it is useful to generate lower and upper bounds for the desired solution. The problems under consideration are weakly nonlinear two point boundary value problems. It is possible to introduce special discretization techniques generating bounds by means of monotonicity principles. More precisely, the proposed approach takes advantage of weak maximum principles leading to operators of monotone kind. In contrast to other approaches it is not necessary to use a priori information or correction terms based on estimations of the local error.

In the present paper we summarize the results from [2], [3], [4], [6] supplemented by some new considerations.

### 2. A first order method

Let us consider the weakly nonlinear boundary value problem

$$(2.1) \quad \begin{aligned} -u'' + g(\cdot, u) &= 0 \quad \text{in } \Omega = (0, 1), \\ u(0) &= u(1) = 0, \end{aligned}$$

under the following assumptions on  $g: \bar{\Omega} \times \mathbb{R} \rightarrow \mathbb{R}$ :

- (2.2) (i)  $g(x, s) \leq g(x, t)$  for all  $x \in \bar{\Omega}$  and  $s \leq t$ ;  
 (ii)  $g$  is a continuous function.

Let us denote by  $U$  the Sobolev space  $H_0^1(0, 1)$ , by  $U^*$  the related dual space and by  $\langle \cdot, \cdot \rangle$  the dual pairing between  $U^*$  and  $U$ .

Introducing the operators  $L, G: U \rightarrow U^*$  defined by

$$\begin{aligned} \langle Lu, v \rangle &:= \int_0^1 u' v' dx, \\ \langle Gu, v \rangle &:= \int_0^1 g(x, u(x)) v(x) dx \end{aligned}$$

[257]

we formulate (2.1) as an operator equation: Find  $u \in U$  such that

$$(2.3) \quad Lu + Gu = 0.$$

Under the proposed assumptions problem (2.3) admits a unique solution, the operator  $L + G$  is strongly monotone and of monotone kind.

Now let be given some grid on the interval  $[0, 1]$ , i.e.

$$0 = x_0 < x_1 < \dots < x_{N-1} < x_N = 1.$$

The corresponding step sizes and subintervals we denote by  $h_i := x_i - x_{i-1}$ ,  $\Omega_i := (x_{i-1}, x_i)$ ,  $i = 1(1)N$ , the mesh width of the grid is characterized by

$$h := \max_{1 \leq i \leq N} h_i.$$

In the sequel we concentrate ourselves on generating upper solutions. We call an operator  $G_h: U \rightarrow U^*$  *bounding operator* if the inequality

$$(2.4) \quad Gv \geq G_h v \quad \text{for all } v \in U$$

holds. It is possible to construct bounding operators using

$$(2.5) \quad G_h := P_h G \quad \text{with } w \geq P_h w \text{ for all } w \in C(\bar{\Omega}).$$

We obtain a first order method by setting, for instance,

$$(2.6) \quad [P_h w](x) := \min_{\xi \in \bar{\Omega}_i} g(\xi, w(\xi)) \quad \text{for } x \in \Omega_i, i = 1(1)N.$$

Our approximate problem is defined by: Find  $u_h \in U$  such that

$$(2.7) \quad Lu_h + G_h u_h = 0.$$

If  $G_h$  is defined by (2.5), (2.6) one can interpret (2.7) as a  $C^1$ -collocation method or as a three point difference scheme. To explain these facts let us introduce the representation

$$(2.8) \quad u_h(x) = \sum_{i=1}^{N-1} u_i \varphi_i(x) + \sum_{i=1}^N w_i \psi_i(x)$$

with ( $u_h$  is piecewise quadratic because  $P_h w$  being piecewise constant)

$$\varphi_i(x) = \begin{cases} (x - x_{i-1})/h_i, & x \in \bar{\Omega}_i, \\ (x_{i+1} - x)/h_{i+1}, & x \in \bar{\Omega}_{i+1}, \\ 0 & \text{otherwise;} \end{cases}$$

$$\psi_i(x) = \begin{cases} \frac{1}{2}(x - x_{i-1})(x_i - x), & x \in \bar{\Omega}_i, \\ 0 & \text{otherwise.} \end{cases}$$

Then, our approximate problem (2.7) is equivalent to the  $C^1$ -collocation method

$$(2.9) \quad u_h \in C^1[0, 1], \quad [Lu + Gu](\xi_i) = 0, \quad i = 1(1)N,$$

where the *collocation points*  $\xi_i$  are defined by

$$g(\xi_i, u_h(\xi_i)) = \min_{\xi \in \bar{\Omega}_i} g(\xi, w(\xi)).$$

The corresponding three-point difference scheme is characterized by

$$(2.10) \quad -D_h u_i + g_h(u_{i-1}, u_i, u_{i+1}) = 0$$

with

$$(2.11) \quad \begin{aligned} \text{(i)} \quad D_h u_i &= \frac{2}{h_i + h_{i+1}} \left( \frac{u_{i-1}}{h_i} - \left( \frac{1}{h_i} + \frac{1}{h_{i+1}} \right) u_i + \frac{u_{i+1}}{h_{i+1}} \right) \\ \text{(ii)} \quad g_h &= -\frac{h_i}{h_i + h_{i+1}} w_i(u_{i-1}, u_i) - \frac{h_{i+1}}{h_i + h_{i+1}} w_{i+1}(u_i, u_{i+1}) \end{aligned}$$

and

$$\begin{aligned} &w_i(u_{i-1}, u_i) \\ &+ g\left(\xi_i, u_{i-1} \frac{x_i - \xi_i}{h_i} + u_i \frac{\xi_i - x_{i-1}}{h_i} + w_i \frac{1}{2} (\xi_i - x_{i-1})(x_i - \xi_i)\right) = 0. \end{aligned}$$

Now, let us assume

$$(2.12) \quad |g(x, s) - g(y, t)| \leq l(r)(|x - y| + |s - t|) \quad \text{for } |s| \leq r, |t| \leq r,$$

where  $l: R^+ \rightarrow R^+$  denotes some nondecreasing function. For the method defined by (2.7) with (2.5), (2.6) we proved in [2]

**THEOREM 1.** *There exists an  $h^* > 0$  such that the discrete problem for all  $h \in (0, h^*]$  admits a solution, further the relations  $u(x) \leq u_h(x)$  and  $\|u - u_h\| \leq Ch$  are valid.*

Here and in the following  $\|\cdot\|$  denotes the maximum norm. The proof of Theorem 1 is based on the strong monotonicity and inverse-monotonicity of the operator  $L + G$ .

The definition of the operator  $P_h$  according to (2.6) is connected with some trouble in the implementation of the method. But, of course, it is possible to define  $P_h$  in some special cases in a more suitable way taking into account special properties of  $g(x, u)$ . Let us assume, for instance, that the function  $g$  admits the representation

$$(2.13) \quad g(x, u) = e(x) + f(u)$$

with some Lipschitz-continuous, monotone function  $f(\cdot)$ .

We denote the finite dimensional space of all functions  $v_h(x)$  of the form

$$(2.14) \quad v_h(x) = \sum_{i=1}^{N-1} v_i \varphi_i(x) + \sum_{i=1}^N w_i \psi_i(x)$$

by  $V_h$  and define a mapping from  $V_h$  into the space of piecewise constant functions using

$$(2.15) \quad [Pv_h](x) := e_i + f(\mu_i) \text{ for all } x \in \Omega_i$$

with

$$e_i \leq e(x) \quad \text{on } \Omega_i, \quad |e(x) - e_i| \leq Ch_i \quad \text{for all } x \in \Omega_i$$

$$\mu_i = \min(v_{i-1}, v_i) + \frac{1}{8} \min(0, w_i) h_i^2.$$

Further we define a mapping  $\pi: C[0, 1] \rightarrow V_h$  by

$$(2.16) \quad [\pi y](x) := \sum_{i=1}^{N-1} v_i \varphi_i(x) + \sum_{i=1}^N w_i \psi_i(x)$$

with

$$v_i = y(x_i),$$

$$w_i = \frac{4}{h_i^2} \left( y(x_{i-1/2}) - \frac{y(x_i) + y(x_{i-1})}{2} \right)$$

( $x_{i-1/2}$  is the midpoint of  $\Omega_i$ ). Finally, a “bounding operator” is defined by

$$(2.17) \quad G_h = P\pi.$$

The new “bounding operator” (2.17) only fulfills

$$(2.18) \quad Gv_h \geq G_h v_h \quad \text{for all } v_h \in V_h,$$

but this property allows us to prove that  $u_h(x)$  is an upper bound of the exact solution because from

$$(u'_h, v') + (Gu_h, v) = ((G - G_h)u_h, v) \geq 0 \quad \text{for all } v \geq 0$$

it follows  $u_h(x) \geq u(x)$ .

Similarly, it is possible to transfer all other parts of the proof of Theorem 1 to our new bounding operator, thus the new discrete problem admits a solution, too. An error estimation is based on

$$\|u - u_h\| \leq C \|Gu_h - g(\cdot, u_h)\|_{L^\infty},$$

and it is not difficult to show that (2.14), (2.15) result in a first order method.

### 3. The application of the first order method to a singular perturbation problem

Of course, it is also possible to analyze our discretization method, described by (2.7) with (2.5), (2.6), in the context of a finite difference method.

We shall do this for the singularly perturbed boundary value problem

$$(3.1) \quad \begin{aligned} -\varepsilon^2 u'' + g(\cdot, u) &= 0, \\ u(0) &= u(1) = 0 \end{aligned}$$

additionally assuming  $(0 < \varepsilon \ll 1)$

$$(3.2) \quad \begin{aligned} (i) \quad &g \in C^2, \\ (ii) \quad &0 < \beta^2 \leq g_u(x, u) \leq B. \end{aligned}$$

It is well known that now the solution  $u$  belongs to  $C^4$  and that near the boundaries boundary layers exist.

For the consistency error we have

$$(3.3) \quad \begin{aligned} |r_h u(x_i)| &\leq K\varepsilon^2 (|h_{i+1} - h_i| |u'''(x_i)| \\ &\quad + \max(h_i^2, h_{i+1}^2) \max_{s \in \bar{\Omega}_i} |u^{(4)}(s)|) \\ &\quad + K \max(h_i, h_{i+1}) \max_{s \in \Omega_i} |u'(s)|. \end{aligned}$$

(we denote by  $K$  some constant which is independent of  $\varepsilon$  and  $h$ ).

In contrast to the usual discretization method

$$(3.4) \quad -\varepsilon^2 D_h u_i + g(x_i, u_i) = 0$$

the difference scheme (2.10), (2.11) applied to our singular perturbation problem is not stable uniformly in  $\varepsilon$ . This is due to the fact that the partial derivatives  $\partial g_u / \partial u_{i-1}$ ,  $\partial g_u / \partial u_{i+1}$  are not negative and thus, in general, the system (2.10) does not correspond to an  $M$ -function. But in the special case

$$(3.5) \quad h \leq \sqrt{2/B} \varepsilon$$

the theory of  $M$ -functions is applicable and the discrete problem is under this restriction uniformly stable with respect to the parameter.

The consistency error is not uniformly bounded in  $\varepsilon$  because the derivatives of the exact solution grow unboundedly near  $x = 0$  and  $x = 1$  as  $\varepsilon$  tends to 0. For the derivatives it holds the estimation

$$(3.6) \quad |u^{(k)}(x)| \leq K(1 + \varepsilon^{-k} (\exp(-\beta x/\varepsilon) + \exp(-\beta(1-x)/\varepsilon))).$$

It is possible to bound the consistency error uniformly by using a special discretization mesh. Let us use the grid generating function due to Vulanovic [12]

$$(3.7) \quad \lambda(t) = \begin{cases} \psi(t) := A\varepsilon t/(q-t), & t \in [0, \tau], \\ \pi(t) := \psi(\tau) + \psi'(\tau)(t-\tau), & t \in [\tau, \frac{1}{2}], \\ 1 - \lambda(1-t), & t \in [\frac{1}{2}, 1], \end{cases}$$

where  $q, A$  are constants with  $q \in (0, \frac{1}{2})$ ,  $A \in (0, q/\varepsilon)$ ,  $\tau$  satisfies

$$(3.8) \quad \tau = (q - (Aq\varepsilon(1 - 2q + 2A\varepsilon))^{1/2}) / (1 + 2A\varepsilon)$$

and define the special grid by

$$(3.9) \quad x_i = \lambda(t_i), \quad t_i = ih^*, \quad i = 0(1)N, \quad N = 2N_0.$$

In [12] it was shown that the consistency error for the approximation of  $u''(x_i)$  admits the order 2 uniformly with respect to  $\varepsilon$  and that the estimate

$$\max(h_i, h_{i+1})\varepsilon^{-1} \leq Kh^*$$

holds. Thus, we have

**THEOREM 2.** *Let us apply the first order enclosing discretization technique described by (2.7) with (2.5), (2.6) to the singularly perturbed boundary value problem (3.1) under the additional assumptions (2.3). Let us use the special grid defined by (3.7), (3.8) and assume that  $\varepsilon$  is not too small, that means,  $h \leq \sqrt{2/B\varepsilon}$ . Then, it holds*

$$|||u - u_h||| \leq Kh^*$$

where  $K$  does not depend on the parameter  $\varepsilon$ .

#### 4. Higher order methods

To generate higher order methods it is necessary to improve the approximation properties of the operator  $G_h$  (respectively  $P_h$ , if  $G_h$  is defined by (2.5)). In [4] we described such bounding operators  $G_h$  defined by (2.5) using piecewise interpolation and shifting to define  $P_h$ .

Let us choose some natural number  $k \geq 1$  and define

$$\mathcal{P}_k = \{v \in L_2(0, 1): v|_{\Omega_i} \text{ polynomial of degree } \leq k\}.$$

We use an equidistant auxiliary grid on every subinterval  $\bar{\Omega}_i$  given by

$$\sigma_j^i := x_{i-1} + \frac{j}{k} h_i, \quad j = 0(1)k,$$

and define a piecewise interpolation operator  $S_k: C[0, 1] \rightarrow \Omega_k$  by  $S_k u|_{\Omega_i} :=$  interpolation polynomial of  $u$  with knots  $\sigma_j^i$  ( $j = 0(1)k$ ). Further we introduce the operator  $p: C[0, 1] \rightarrow \mathcal{P}_0$  by

$$(4.1) \quad [pw](x) := \min_{\xi \in \bar{\Omega}_i} w(\xi) \quad \text{for all } x \in \Omega_i$$

and define

$$(4.2) \quad P_h := S_k + p(I - S_k)$$

( $I$  is the identity).

It is not so easy to prove that the generated discrete problem (2.7) admits a solution. In our proof we used the theory of pseudomonotone operators [4] and an auxiliary variational inequality [5], respectively. The convergence proof works like the convergence proof for finite element methods. Analogously as for collocation methods, we obtain the following result:

**THEOREM 3.** *There exist an  $h^* > 0$  such that the discrete problem for all  $h \in (0, h^*]$  admits a solution. If all partial derivatives of order less or equal to  $k+1$  are continuous on  $\Omega_i$  for  $i = 1(1)N$  then it holds the error estimate*

$$\|u - u_h\| \leq Ch^{k+1}.$$

The numerical experiments meet the theoretical expectation [4].

### 5. Grid generation via enclosing discretization

In this section we propose a principle for the generation of grids using information from the lower and upper solution. The basic idea is to subdivide these intervals where the difference between the upper and lower solutions is relatively large.

Let some initial grid  $Z^1 = \{x_i^1: i = 0, \dots, N_1\}$  be given. The related subintervals we denote by  $\Omega_i^1$  and assume  $h(Z^1)$  to be small enough that the discrete problems possess a solution. The algorithm under consideration generates a sequence  $(Z^k)$  of grids. The bounding operators on the grids  $Z^k$  we denote by  $G^k$ .

#### *Algorithm*

**Step 1.** Let some initial grid  $Z^1$  be given. Select some  $\varrho \in (0, 1)$  and set  $k = 1$ .

**Step 2.** Determine  $\underline{u}^k(x)$ ,  $\bar{u}^k(x)$  by solving the corresponding discrete problems (2.7) (with  $\underline{G}^k$ ,  $\bar{G}^k$  instead of  $G_h$ ).

**Step 3.** With the notation

$$d_i^k := \max_{x \in \Omega_i^k} (\bar{u}^k(x) - \underline{u}^k(x)) \quad (i = 1, \dots, N_k),$$

$$D^k := \max_{1 \leq i \leq N_k} d_i^k,$$

$$I_k := \{i \in \{1, \dots, N_k\}: d_i^k \geq \varrho D^k\},$$

$$y_i^k := \frac{1}{2}(x_{i-1}^k + x_i^k) \quad (i \in I_k)$$

define a new grid by

$$Z^{k+1} = Z^k \cup \{y_i^k: i \in I_k\}.$$

Set

$$N_{k+1} = N_k + \text{card } I_k$$

and denote the grid points contained in  $Z^{k+1}$  by  $x_i^{k+1}$  ( $i = 0, 1, \dots, N_{k+1}$ ).  
Reset

$$k^{\text{new}} := k^{\text{old}} + 1$$

and go to step 2.

In [3] we succeeded in proving the following

**THEOREM 4.** *Let us suppose that in step 2 of the algorithm above the first order method characterized by the bounding operator defined by (2.5), (2.6) (and the corresponding operator for lower solutions) is used. Then, the functions  $\underline{u}^k(x)$ ,  $\bar{u}^k(x)$  generated have the properties*

$$\underline{u}^k(x) \leq \underline{u}^{k+1}(x) \leq u(x) \leq \bar{u}^{k+1}(x) \leq \bar{u}^k(x) \quad \text{for all } x \in [0, 1],$$

$$\lim_{k \rightarrow \infty} \|\bar{u}^k(x) - \underline{u}^k(x)\| = 0.$$

In numerical experiments we obtained improved error bounds using the grid generation algorithm in connection with higher order methods in step 2, too. However, the theoretical foundation of the feedback grid generation cannot be carried out similarly to the first-order technique. The proof of Theorem 4 is based on the following two statements. First, for the first-order method it holds: From  $v, w \in C^1[0, 1]$  and

$$\begin{aligned} Lv + G_h v &\leq Lw + G_h w, \\ v(0) &\leq w(0), \quad v(1) \leq w(1), \end{aligned}$$

it follows  $v(x) \leq w(x)$  for sufficiently small  $h$ . And, second, for the first order method an enclosure  $\underline{u}_Z \leq u \leq \bar{u}_Z$  is sharpened using a finer grid. It is an open problem whether or not higher order methods preserve these properties.

## 6. A second approach to enclosing discretizations

As in Section 2 we start from the formulation of problem (2.1) as an operator equation: Find some  $u \in U$  such that

$$(6.1) \quad Lu + Gu = 0.$$

In contrast to the discretization technique described above we now shall formulate an iteration process of the form

$$(6.2) \quad \begin{aligned} \text{(i)} \quad & (L + D)u^{k+1} = s^k; \\ \text{(ii)} \quad & s^{k+1} = \min(s^k, Qu^{k+1}) \end{aligned}$$

using some auxiliary operators  $D, Q$  such that  $\{u^k\}$  converges monotone to an upper (lower) solution of the original problem (6.1). This iteration process is similar to the technique of Sattinger [10]

$$(L + \varrho I)\hat{u}^{k+1} = (\varrho I - G)\hat{u}^k$$



generating a monotone sequence but we propose a modification which is numerically constructive.

First, we define the operator  $D$  by

$$(6.3) \quad [Dw](x) := d_i w(x) \quad \text{for all } x \in \Omega_i$$

where the nonnegative constants  $d_i$  are to specify in such a manner that the required property of  $D - G$  is fulfilled. Now we choose some finite-dimensional space  $W_h$  of the type

$$(6.4) \quad W_h = \{w \in L_2(0, 1): w|_{\Omega_i} \text{ polynomial of degree } \leq l\}$$

and define

$$(6.5) \quad V_h = (L + D)^{-1} W_h.$$

We choose  $s^0 \in W_h$ ,  $u^0 \in V_h$  and define  $Q$  to be a mapping  $Q: V_h \rightarrow W_h$  which satisfies

$$(6.6) \quad Qw \geq (D - G)w \quad \text{for all } w \in V_h.$$

More precisely, we need initial values  $\underline{s}^0$ ,  $\underline{u}^0$  and  $\bar{s}^0$ ,  $\bar{u}^0$  for the iteration process generating lower and upper solutions, respectively. Let us suppose

$$(6.7) \quad \begin{aligned} & \text{(i)} \quad \underline{u}^0 \leq \bar{u}^0; \\ & \text{(ii)} \quad (D - G)w_1 \leq (D - G)w_2 \text{ for all } w_{1,2} \text{ with } \underline{u}^0 \leq w_1 \leq w_2 \leq \bar{u}^0; \\ & \text{(iii)} \quad (D - G)\bar{u}^0 \leq \bar{s}^0, (D - G)\underline{u}^0 \geq \underline{s}^0; \\ & \text{(iv)} \quad (L + D)\bar{u}^0 \geq \bar{s}^0, (L + D)\underline{u}^0 \leq \underline{s}^0. \end{aligned}$$

As above we only state the following theorem from [6] concerning the sequence  $\{\bar{u}_k\}$  which tends to an upper solution. We omit the bar in the notation again.

**THEOREM 5.** *Under the assumptions (6.6), (6.7) the iteration process (6.2) generates a sequence  $\{u_k\}$  with the following properties:*

- (i)  $u_{k+1} \leq u_k \leq \dots \leq u_0$ ;
- (ii)  $\{u_k\}$  tends to some  $\{\tilde{u}\}$  in  $H^1(0, 1)$  and  $C[0, 1]$ ;
- (iii) For the limit function  $\tilde{u}$  it holds  $\tilde{u} \geq u$  and

$$(6.8) \quad \|\tilde{u} - u\| \leq C \|Q\tilde{u} - (D - G)\tilde{u}\|_{L^\infty}.$$

The realization of this second approach to enclosing discretizations strongly depends on the concrete definition of the space  $W_h$  and the operator  $Q$ . One expects a first order method in choosing  $W_h$  to be the space of piecewise constant functions, a second order method by choosing  $W_h$  as the space of piecewise linear functions and defining  $Q$  in an adequate way and so on.

In the first case it is obvious to define

$$(6.9) \quad [Qw](x) := \max_{\xi \in \bar{\Omega}_i} (d_i w(\xi) - g(\xi, w(\xi))) \quad \text{for all } x \in \Omega_i.$$

Similarly as in Section 2 it is possible to make the method more practicable under special assumptions on  $g$ . Let us assume again

$$(6.10) \quad g(x, u) = e(x) + f(u)$$

(compare (2.13) and the following statements). Now we have the advantage that it is only necessary to define the operator  $Q$  on  $V_h$ . All functions from  $V_h$  admit a representation

$$(6.11) \quad v_h(x) = \sum_{i=1}^{N-1} v_i \varphi_i^*(x) + \sum_{i=1}^N w_i \psi_i^*(x)$$

with

$$(6.12) \quad \begin{aligned} \text{(i)} \quad \varphi_i^*(x) &= \begin{cases} \sinh(\sqrt{d_i}(x-x_{i-1}))/\sinh(\sqrt{d_i}h_i) & \text{for } x \in \bar{\Omega}_i, \\ \sinh(\sqrt{d_{i+1}}(x_{i+1}-x_i))/\sinh(\sqrt{d_{i+1}}h_{i+1}) & \text{for } x \in \bar{\Omega}_{i+1}, \\ 0 & \text{otherwise;} \end{cases} \\ \text{(ii)} \quad \varphi_i^*(x) &= \begin{cases} (1-\varphi_{i-1}^*(x)-\varphi_i^*(x))/d_i & \text{for } x \in \bar{\Omega}_i, \\ 0 & \text{otherwise.} \end{cases} \end{aligned}$$

Then, it is adequate to define

$$(6.13) \quad [Qv_h](x) := -e_i + d_i \mu_i^* f(\mu_i^*) \quad \text{for all } x \in \Omega_i, v_h \in V_h$$

with

$$\mu_i^* := \max(v_{i-1}, v_i) + \frac{1}{8} \max(0, w_i) h_i^2.$$

The proof of Theorem 5 shows the validity of the estimate

$$(6.14) \quad s^k \geq (D-G) u^{k+1}.$$

Therefore, we could formulate the iteration process (6.2) with (6.9) in the form

$$(6.15) \quad \begin{aligned} \text{(i)} \quad (L+D) u^{k+1} &= s^k; \\ \text{(ii)} \quad s_i^{k+1} &= s_i^k - \min_{\xi \in \bar{\Omega}_i} (s_i^k - (d_i u^{k+1}(\xi) - g(\xi, u^{k+1}(\xi)))). \end{aligned}$$

( $s_i^k$  denotes the restriction of  $s^k$  to  $\Omega_i$ ).

For generating a second order method we choose  $W_h$  as the space of piecewise linear functions and use again the idea of the reformulation of the iteration process in the form (6.15). Let us define  $\xi_i$  by

$$\begin{aligned} s_i^k(\xi_i) - [d_i u^{k+1}(\xi_i) - g(\xi_i, u^{k+1}(\xi_i))] \\ = \min_{\xi \in \bar{\Omega}_i} \{s_i^k(\xi) - [d_i u^{k+1}(\xi) - g(\xi, u^{k+1}(\xi))]\}. \end{aligned}$$

Further we denote by  $t_u^{i,1}(x)$  the linear interpolant of  $d_i u - g(\cdot, u(\cdot))$  with respect to the gridpoints  $x_{i-1}$  and  $x_i$  and by  $t_u^{i,2}(x)$  the linear function on

$[x_{i-1}, x_i]$  which passes  $d_i u - g(\cdot, u(\cdot))$  in  $x_{i-1}$  or  $x_i$  (if  $\xi_i = x_{i-1}$  or  $\xi_i = x_i$ ), majorizes  $d_i u - g(\cdot, u(\cdot))$  over the interval  $[x_{i-1}, x_i]$  and is tangent to  $d_i u - g(\cdot, u(\cdot))$  in some point. Using

$$t_u^i(x) = \begin{cases} t_u^{i,1}(x) & \text{if } t_u^{i,1}(x) \text{ majorizes } d_i u - g(\cdot, u(\cdot)) \\ & \text{over the interval } [x_{i-1}, x_i], \\ t_u^{i,2}(x) & \text{otherwise,} \end{cases}$$

we define the iteration process

$$(6.16) \quad \begin{aligned} & \text{(i) } (L+D)u^{k+1} = s^k \\ & \text{(ii) } s_u^{k+1}(x) = \begin{cases} s_i^k(x) - \min_{\xi \in \bar{\Omega}_i} [s_i^k(x) - (d_i u^{k+1}(\xi) - g(\xi, u^{k+1}(\xi)))] & \text{if } \xi_i \in \Omega_i, \\ t_{u^{k+1}}^i(x) & \text{if } \xi_i = x_{i-1} \text{ or } \xi_i = x_i. \end{cases} \end{aligned}$$

Now, again Theorem 5 holds and from (6.8) we obtain the second order accuracy of the approximation of  $\tilde{u}$  to the exact solution with respect to  $h$ . A more practicable version of a second order method is described in [6].

## 7. Concluding remarks

In our paper we did not speak about the numerical solution of the auxiliary problems generated by the enclosing discretization principle of Section 2 and 4. Special iteration methods have been developed in [5].

The extension of enclosing discretization techniques to partial differential equations leads to significant troubles because no explicit representation of  $L^{-1}W_h$  for the function spaces  $W_h$  used above is available in the multidimensional case.

Some progress has been achieved in extending our enclosing discretization methods to  $1-D$  linear parabolic problems [7], [9] using a modification of the Rothe-method for the discretization in time.

## References

- [1] E. Adams, H. Spreuer, *Konvergente numerische Schrankenkonstruktionen mit Spline-Funktionen für nichtlineare gewöhnliche bzw. lineare parabolische Randwertaufgaben*, K. Nickel (ed.), Springer-Verlag, Berlin-Heidelberg-New York 1975.
- [2] Ch. Großmann, M. Krätzschmar, H.-G. Roos, *Gleichmäßig einschließende Diskretisierungsverfahren für schwach nichtlineare Randwertaufgaben*, Numer. Math. 49 (1986), 95-110.
- [3] Ch. Großmann, H.-G. Roos, *Feedback grid generation via monotone discretization for two-point boundary-value problems*, IMA J. Numer. Anal. 6 (1986), 421-432.
- [4] Ch. Großmann, M. Krätzschmar, H.-G. Roos, *Uniformly enclosing discretization methods of high order for boundary value problems*, Math. Comp. (to appear).

- [5] Ch. Großmann, *Monotone discretization of two-point boundary value problems and related numerical methods*, in Adams, Ansorge, Großmann, Roos (eds.), *Discretization of Differential Equations and Enclosures*, Akademie-Verlag, Berlin 1987.
- [6] Ch. Großmann, *Implementable monotone discretization algorithms* (in preparation).
- [7] G. Koepppe, H.-G. Roos, L. Tobiska, *An enclosure generating modification of the method of discretization in time*, *Comment. Math. Univ. Carolin.* 28 (1987), 447–453.
- [8] K. Nickel, *The construction of a priori bounds for the solution of two-point boundary value problems with finite elements*, *Computing* 23 (1979), 247–265.
- [9] H.-G. Roos, *The Rothe method and monotone discretization for parabolic equations*, in Adams, Ansorge, Großmann, Roos (eds.), *Discretization of Differential Equations and Enclosures*, Akademie-Verlag, Berlin 1987.
- [10] D. H. Sattinger, *Monotone methods in nonlinear elliptic and parabolic boundary value problems*, *Indian J. Math.* 31 (1972), 979–1000.
- [11] H. Spreuer, *A method for the computation of bounds with convergence of arbitrary order for ordinary linear boundary value problems*, *J. Math. Anal. Appl.* 81 (1981), 99–133.
- [12] R. Vulcanovic, *Mesh construction for discretization of singularly perturbed boundary value problems*, *Doctoral Diss.*, Novi Sad 1986.
- [13] P. Wildenauer, *A new method for automatical computation of error bounds for the set of all solutions of nonlinear boundary value problems*, *Computing* 34 (1985), 131–154.

*Presented to the Semester  
Numerical Analysis and Mathematical Modelling  
February 25 – May 29, 1987*

---