

ON REDUCTIVE AND PROJECTIVE CLASSES

BY

KLAUS KAISER (HOUSTON, TEXAS)

The notion of a projective class was introduced by Mal'cev [5]. Let \mathfrak{A} and \mathfrak{B} be axiomatic classes and let K be a set of first order sentences formulated in a language which contains the languages L and M for \mathfrak{A} and \mathfrak{B} , respectively. K describes an axiomatic relation Γ between the models $A \in \mathfrak{A}$ and $B \in \mathfrak{B}$. The models of K are all disjoint unions $A \dot{\cup} B$ which admit an interpretation of the set T of relational symbols occurring in K and not belonging to L or M such that the enlarged structure $(A \dot{\cup} B)^\wedge$ becomes a model of K . The domain of the axiomatic relation Γ is called a *projective subclass* $\mathcal{P} = \mathfrak{A} \Gamma \mathfrak{B}$ of \mathfrak{A} : $A \in \mathcal{P}$ iff $(A \dot{\cup} B)^\wedge \models K$ for some $B \in \mathfrak{B}$ and an interpretation of T iff A acts on B in the sense of Γ (see [1] and [2]).

A class \mathfrak{A} of relational structures of type τ is called *pseudo-axiomatic* or a PC_τ -class if it is the class of τ -reducts of an axiomatic class of extended type $\varrho \supseteq \tau$. This notion is due to Tarski. We use here the term *reductive* instead of pseudo-axiomatic. Clearly, reductive classes are projective.

In this note we first point out that a projective class \mathcal{P} contains a "big" reductive class. But a projective class is also the intersection of reductive classes containing it. The latter statement is a consequence of Keisler's ingenious limit ultraproduct theorem.

THEOREM 1. *Let \mathcal{P} be a projective subclass of the axiomatic class \mathfrak{A} and assume that \mathcal{P} contains only infinite models. Then \mathcal{P} contains a reductive class \mathfrak{R} such that any $A \in \mathcal{P}$ can be elementarily embedded into a model $A^\wedge \in \mathfrak{R}$.*

Proof. Let the projective class \mathcal{P} be defined by the axiomatic relation Γ between \mathfrak{A} and the axiomatic class \mathfrak{B} . We select in \mathcal{P} the class \mathfrak{R} of all $A \in \mathcal{P}$ which admit a realization of Γ by a model $B \in \mathfrak{B}$ with $\text{card}(A) \geq \text{card}(B)$. We first show that \mathfrak{R} is reductive. We extend for this purpose the language L_τ for \mathfrak{A} by the set of relational symbols in M , by the set T of the connecting relations for Γ and by one unary relational symbol R . Let the set K^\wedge say that a τ -reduct A of a model A^\wedge of K^\wedge is a model in \mathfrak{A} ,

the subset defined by R is a model B in \mathfrak{B} and that A acts on B in the sense of Γ . Then $A \in \text{Md } K^\wedge$ iff $A \in \mathfrak{R}$.

For $A_0 \in \mathcal{P}$, we have a $B_0 \in \mathfrak{B}$ such that $(A_0 \dot{\cup} B_0)^\wedge \models K$. We can find an elementary extension $(A_1 \dot{\cup} B_1)^\wedge$ of $(A_0 \dot{\cup} B_0)^\wedge$ together with an injective map $f_0: B_0 \rightarrow A_1$ such that $A_1 - \text{im}(f_0)$ is infinite. There is an elementary extension $(A_2 \dot{\cup} B_2)^\wedge$ of $(A_1 \dot{\cup} B_1)^\wedge$ together with an injective map $f_1: B_1 \rightarrow A_2$, where f_1 extends f_0 and where $A_2 - \text{im}(f_1)$ is infinite. Continuing in this way, we get an elementary chain of extensions and a chain of injections:

$$\begin{array}{c} (A_0 \dot{\cup} B_0)^\wedge \hookrightarrow (A_1 \dot{\cup} B_1)^\wedge \hookrightarrow (A_2 \dot{\cup} B_2)^\wedge \hookrightarrow \dots \\ \downarrow \quad \downarrow \quad \downarrow \\ f_0 \quad f_1 \quad f_2 \\ \downarrow \quad \downarrow \quad \downarrow \\ f_0 \subset f_1 \subset \dots \end{array}$$

By the well-known Tarski-Vaught theorem, the union A_ω of the A_i acts on the union B_ω of the B_i in the sense of Γ and, moreover, we have an injection

$$f_\omega: B_\omega \rightarrow A_\omega, \quad f_\omega = \bigcup_{i=1}^{\infty} f_i.$$

If \mathfrak{B} contains only infinite models, then we can interchange in the last proof the roles of A and B successively, to get an elementary embedding

$$(A_0 \dot{\cup} B_0)^\wedge \hookrightarrow (\overline{A} \dot{\cup} \overline{B})^\wedge, \quad \text{where } \text{card}(\overline{A}) = \text{card}(\overline{B}).$$

Call a realization of Γ *balanced* if A acts on itself, $A \Gamma B$, the set B a double of A . Then we have the following

COROLLARY. *Let Γ be an axiomatic relation between the axiomatic classes \mathfrak{A} and \mathfrak{B} , where \mathfrak{A} and \mathfrak{B} contain only infinite models. Let \mathfrak{R} be the reductive subclass of $\mathcal{P} = \mathfrak{A} \Gamma \mathfrak{B}$ given by all models which admit a balanced realization of Γ . Then any $A \in \mathfrak{A}$ can be elementarily embedded into a model $\overline{A} \in \mathfrak{R}$.*

By Theorem 1, any projective class of infinite models contains a reductive class dense with respect to the axiomatic closure. Denote by $\mathcal{C}_{\mathcal{P}}$ the topological closure defined by the projective classes [3], and by $\mathcal{C}_{\mathfrak{R}}$ the topological closure given by the reductive classes.

THEOREM 2 (Keisler [4]). *Let \mathfrak{A} be a class of relational structures of type τ . Then $A \in \mathcal{C}_{\mathfrak{R}} \mathfrak{A}$ iff A is isomorphic to a limit ultraproduct with factors in \mathfrak{A} .*

-Recall that, for a family $(A_i)_{i \in I}$, ultrafilter u over $I \times J$ and filter \mathcal{G} over $I \times J \times J$, the limit ultraproduct

$$A = \prod_{(i,j) \in I \times J} A_{ij} / u | \mathcal{G}, \quad A_{ij} = A_i,$$

is defined as the subset of all equivalence classes a of $\prod_{i,j} A_{ij}/u$ which contain a representative f such that

$$\text{eq}(f) = \{(i, j, j') \mid f(i, j) = f(i, j')\} \in \mathfrak{G}.$$

THEOREM 3. *A projective class \mathcal{P} is closed under limit ultraproducts.*

Proof (cf. [1]). Let $(A_i)_{i \in I}$ be a family of models in \mathcal{P} , where \mathcal{P} is a projective class defined by the relation Γ between the axiomatic class \mathfrak{A} and the axiomatic class \mathfrak{B} . We then have $B_i \in \mathfrak{B}$ such that $(A_i \dot{\cup} B_i)^\wedge \models K$, K describing Γ . It is easy to see that there is a canonically given isomorphism

$$\prod_{i,j} (A_{ij} \dot{\cup} B_{ij})/u \mid \mathfrak{G} \cong \prod_{i,j} A_{ij}/u \mid \mathfrak{G} \dot{\cup} \prod_{i,j} B_{ij}/u \mid \mathfrak{G},$$

and this isomorphism can be made to one of the extended structures, showing that the limit ultraproduct

$$A = \prod_{i,j} A_{ij}/u \mid \mathfrak{G} \in \mathfrak{A}$$

acts on

$$B = \prod_{i,j} B_{ij}/u \mid \mathfrak{G} \in \mathfrak{B}$$

in the sense of Γ .

By Keisler's theorem we now have the following

COROLLARY. *Let \mathcal{P} be a projective class and let $A \notin \mathcal{P}$. Then there is a reductive class $\mathfrak{R} \supseteq \mathcal{P}$ with $A \notin \mathfrak{R}$, $\mathcal{C}_{\mathfrak{R}} = \mathcal{C}_{\mathcal{P}}$.*

In [3] we proved that $A' \in \mathcal{C}_{\mathcal{P}}\mathfrak{A}$ iff there is an ultraproduct

$$A = \prod_{i \in I} A_i/u, \quad A_i \in \mathfrak{A},$$

such that any projective class containing not only A , but also almost all factors, contains also A' . If equation (*) in Keisler [4], p. 219, holds for \mathfrak{A} , then one can dispense with the restriction for \mathcal{P} to contain almost all factors; but the truth of (*), in general, is unsettled.

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