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LAWLER AND BELL'S METHOD OF DISCRETE OPTIMIZATION

1. Procedure declaration. The procedure *doptLB* solves the following problem:

Find the minimum of $g_0(x)$ under the restrictions

$$g_{11}(x) - g_{12}(x) \geq 0, \quad g_{21}(x) - g_{22}(x) \geq 0, \quad \dots, \quad g_{m1}(x) - g_{m2}(x) \geq 0,$$

where $x = (x_1, x_2, \dots, x_n)$, $x_j = 0, 1$, and the functions $g_0, g_{11}, g_{12}, g_{21}, g_{22}, \dots, g_{m1}, g_{m2}$ are monotonically non-decreasing with respect to every x_j .

Data:

n — number of zero-one variables x_j ,

m — number of restrictions $g_{i1}(x) - g_{i2}(x) \geq 0$,

f — function with procedure heading: **real procedure** $f(k, x)$; **integer** k ; **integer array** x ; which for $k = 0$ calculates $g_0(x)$, for $k = 1, 2, \dots, m$ calculates $g_{k1}(x)$, and for $k = m + 1, m + 2, \dots, 2m$ calculates $g_{k-m,2}(x)$, given x in $x[1:n]$,

max — maximum allowable number of type **real**.

Results:

$x[1:n]$ — the solution vector,

z — optimum value of g_0 ; if $z = max$ on exit, then there is no feasible solution to the problem.

2. Method used. The algorithm was coded according to the method given by Lawler and Bell in [1].

3. Certification. Several examples, including those given in [1], were solved and correct results obtained. The calculations were done on the ODRA 1204 computer.

```

procedure doptLB(n,m,f,x,z,max);
  value n,m;
  integer n,m;
  real z,max;
  integer array x;
  real procedure f;
  begin
    integer i;
    real zg,gr;
    integer array xx,xg[1:n];
    array g[1:m];
    for i:=1 step 1 until n do
      xx[i]:=0;
    for i:=1 step 1 until m do
      if f(i,xx)<f(m+i,xx)
        then go to ntrivsol;
    z:=f(0,xx);
    for i:=1 step 1 until n do
      x[i]:=xx[i];
    go to finish;
  ntrivsol:
    xx[n]:=1;
    z:=max;
  repeat:
    zg:=f(0,xx);
    for i:=1 step 1 until n do
      xg[i]:=xx[i];
    for i:=n step -1 until 1 do
      if xx[i]=0
        then xg[i]:=1

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    else go to END;
END:
    if  $z_g > z$ 
        then go to skip;
    for  $i:=1$  step 1 until  $m$  do
        begin
             $g_r := g[i] := f(m+i, x_x)$ ;
            if  $f(i, x_g) < g_r$ 
                then go to skip
            end  $i$ ;
        for  $i:=1$  step 1 until  $m$  do
            if  $f(i, x_x) < g[i]$ 
                then go to next;
         $z := z_g$ ;
        for  $i:=1$  step 1 until  $n$  do
             $x[i] := x_x[i]$ ;
    skip:
        for  $i:=1$  step 1 until  $n$  do
             $x_x[i] := x_g[i]$ ;
    next:
        for  $i:=n$  step -1 until 1 do
            if  $x_x[i] = 1$ 
                then  $x_x[i] := 0$ 
            else
                begin
                     $x_x[i] := 1$ ;
                    go to repeat
                end  $i, x_x[i] = 0$ ;
    finish:
        end doptLB

```

Reference

- [1] E. L. Lawler and M. D. Bell, *A method for solving discrete optimization problems*, Operat. Res. 14 (1966), p. 1098-1112.

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ALGORYTM 23

METODA LAWLERA I BELLA OPTYMALIZACJI DYSKRETNEJ

STRESZCZENIE

Procedura *doptLB* rozwiązuje następujący problem:
Znaleźć minimum funkcji $g_0(x)$ dla warunków

$$g_{11}(x) - g_{12}(x) \geq 0, \quad g_{21}(x) - g_{22}(x) \geq 0, \quad \dots, \quad g_{m1}(x) - g_{m2}(x) \geq 0,$$

gdzie $x = (x_1, x_2, \dots, x_n)$, $x_j = 0, 1$, oraz funkcje $g_0, g_{11}, g_{12}, g_{21}, g_{22}, \dots, g_{m1}, g_{m2}$ są monotonicznie niemalejące ze względu na każdą ze zmiennych x_j .

Dane:

n – liczba zmiennych zero-jedynkowych x_j ,

m – liczba ograniczeń $g_{i1}(x) - g_{i2}(x) \geq 0$,

f – funkcja o nagłówku: **real procedure** $f(k, x)$; **integer** k ; **integer array** x ; która dla $k = 0$ oblicza $g_0(x)$, dla $k = 1, 2, \dots, m$ oblicza $g_{k1}(x)$ i dla $k = m + 1, m + 2, \dots, 2m$ oblicza $g_{k-m,2}(x)$, gdzie x jest dane w $x[1:n]$,

max – największa dopuszczalna liczba typu rzeczywistego.

Wyniki:

$x[1:n]$ – wektor rozwiązania,

z – optymalna wartość funkcji $g_0(x)$; jeżeli po wyjściu z procedury $z = max$, to problem nie ma rozwiązania dopuszczalnego.

Algorytm oparty jest na metodzie podanej w [1]. Poprawność procedury sprawdzono na wielu przykładach na m. c. Odra 1204.